

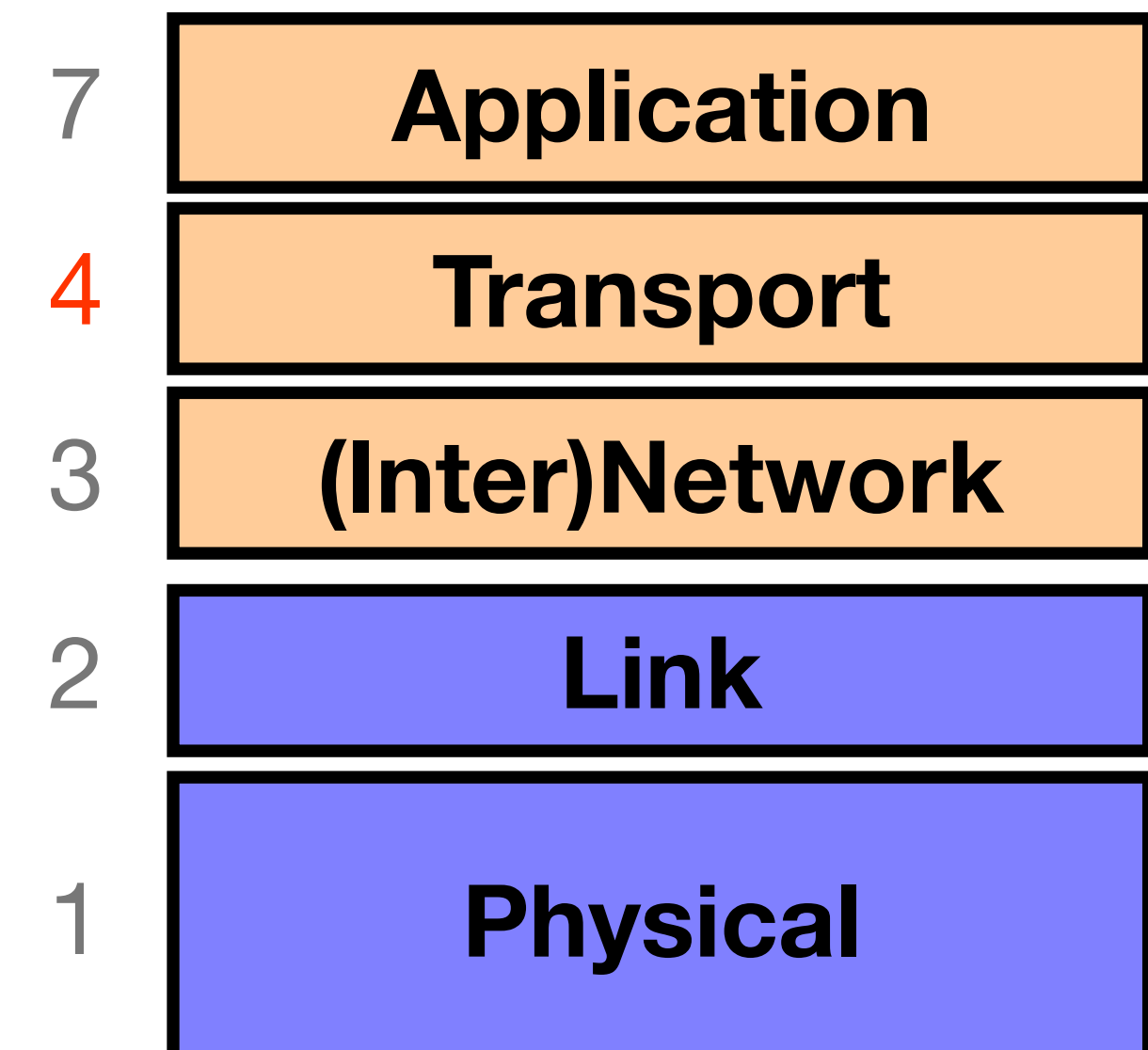
# Lecture 18: Network Security

# Announcements

- Project 2 design doc due Friday
- Networking tutorial, Saturday 3/7, 5-7pm, in HP Auditorium (306 Soda)

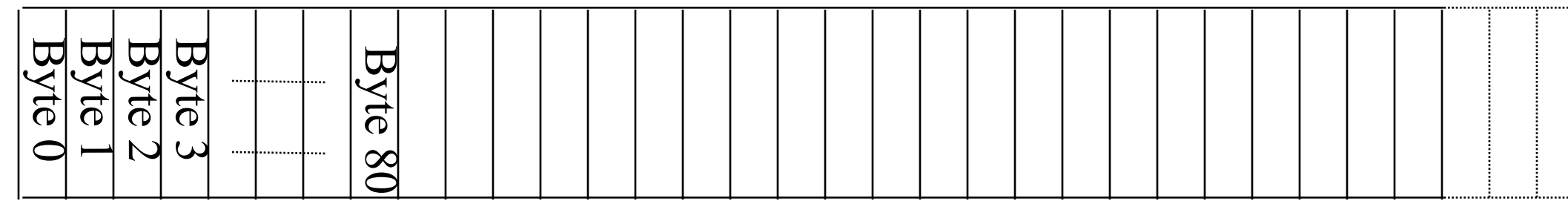
# “Best Effort” is Lame! What to do?

- #1 workhorse: TCP (Transmission Control Protocol)
- Service provided by TCP:
  - Connection oriented (explicit set-up / tear-down)
    - End hosts (processes) can have multiple concurrent long-lived communication
  - **Reliable**, in-order, *byte-stream* delivery
    - Robust detection & retransmission of lost data



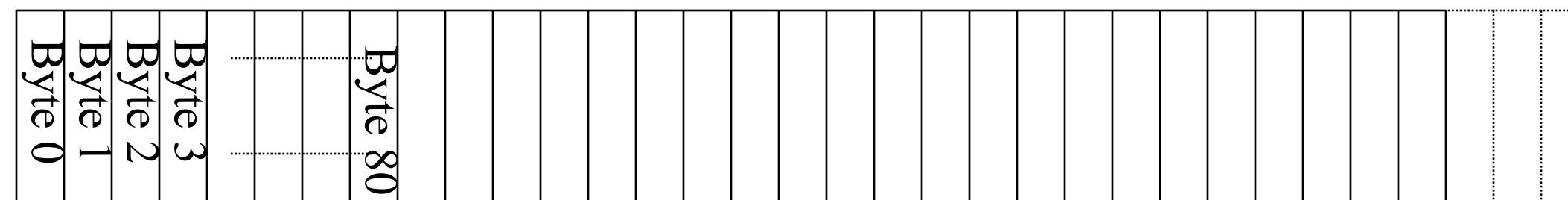
# TCP “Bytestream” Service

# Process A on host H1



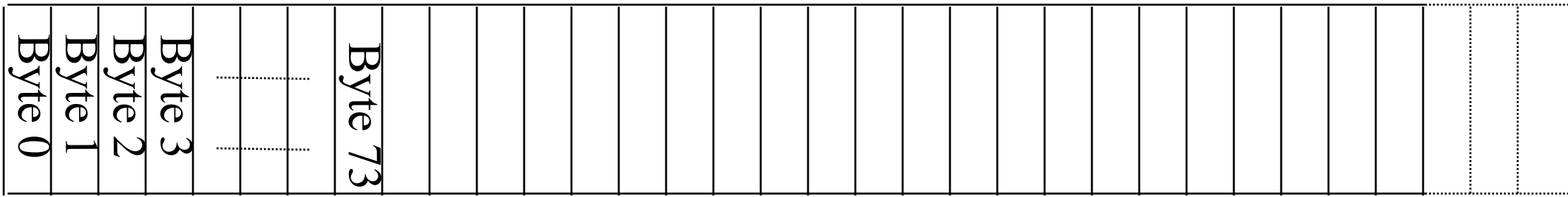
Hosts don't ever see packet boundaries, lost or corrupted packets, retransmissions, etc.

Process B  
on host H2



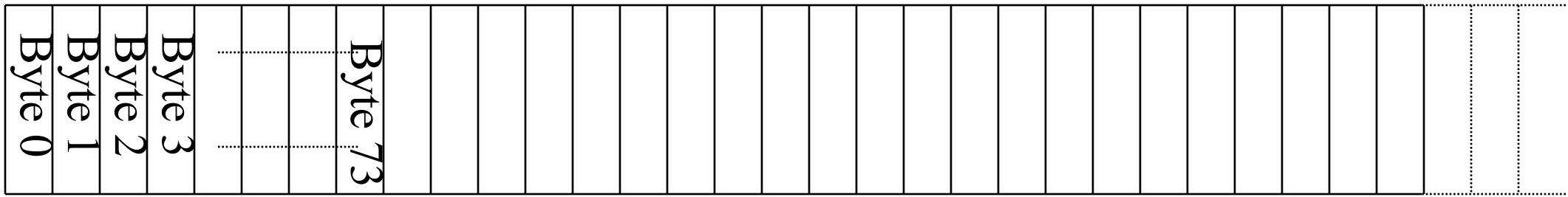
# Bidirectional communication:

Process B on host H2

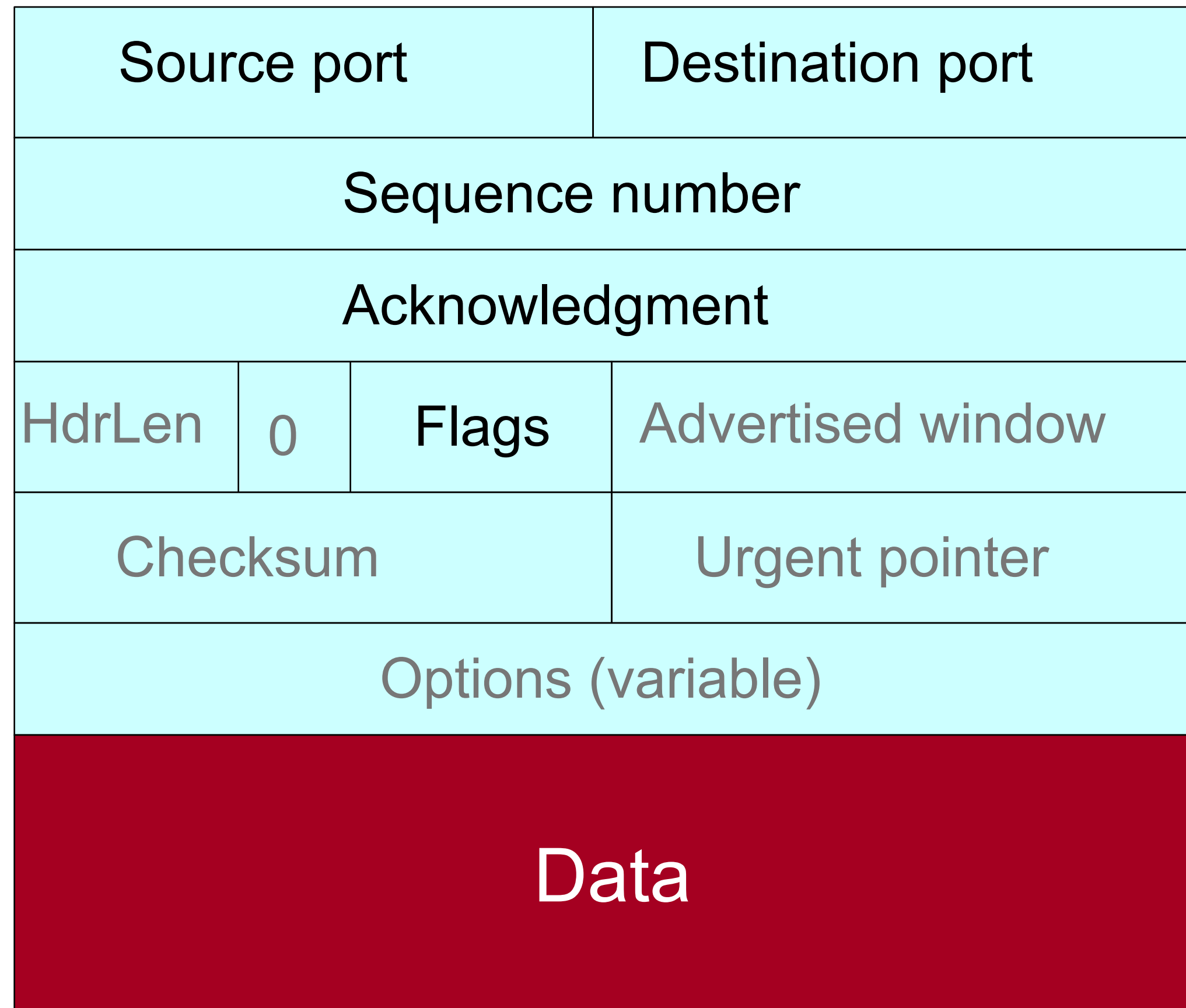


There are two separate bytestreams, one in each direction

Process A  
on host H1

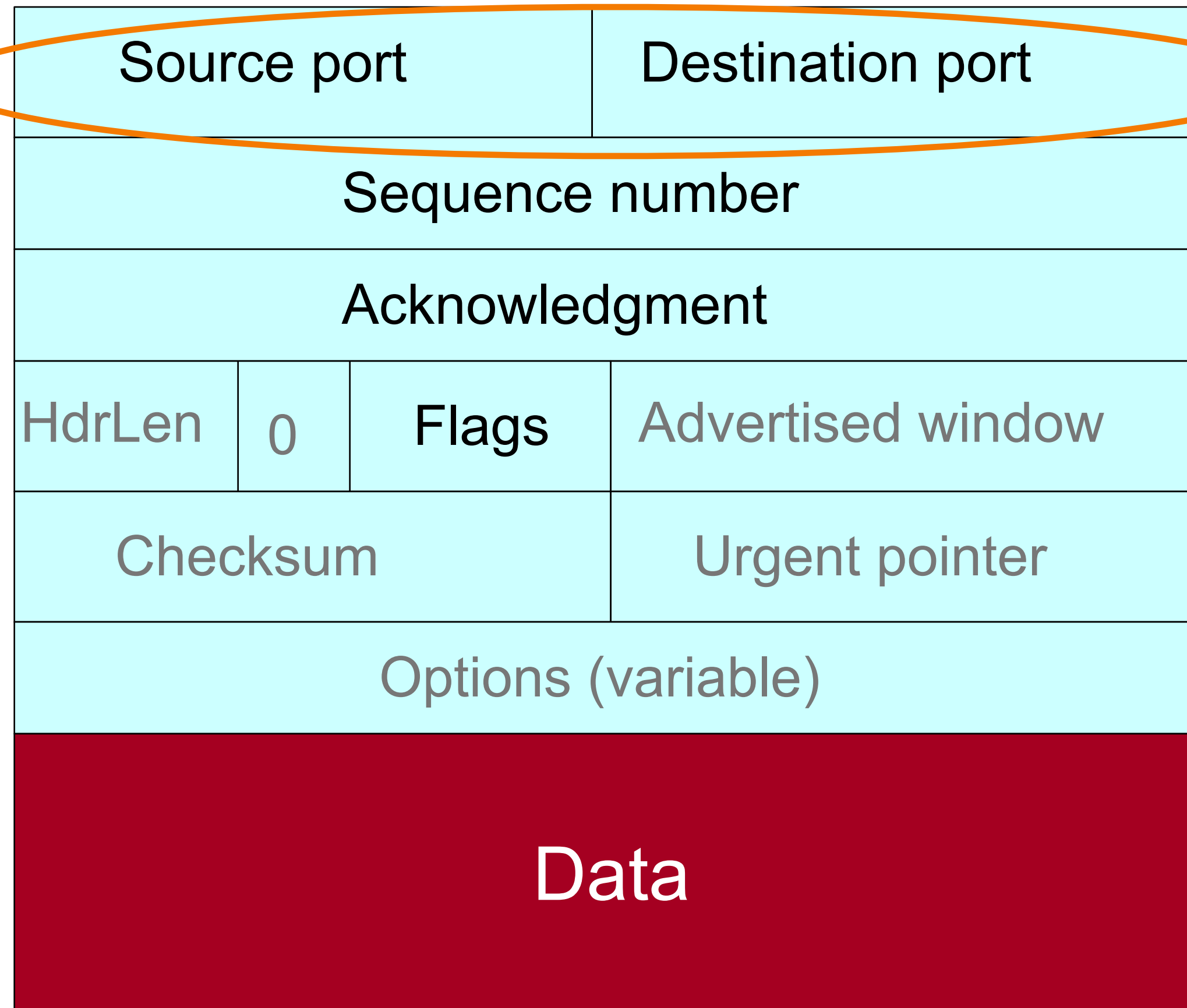


# TCP Header



# TCP Header

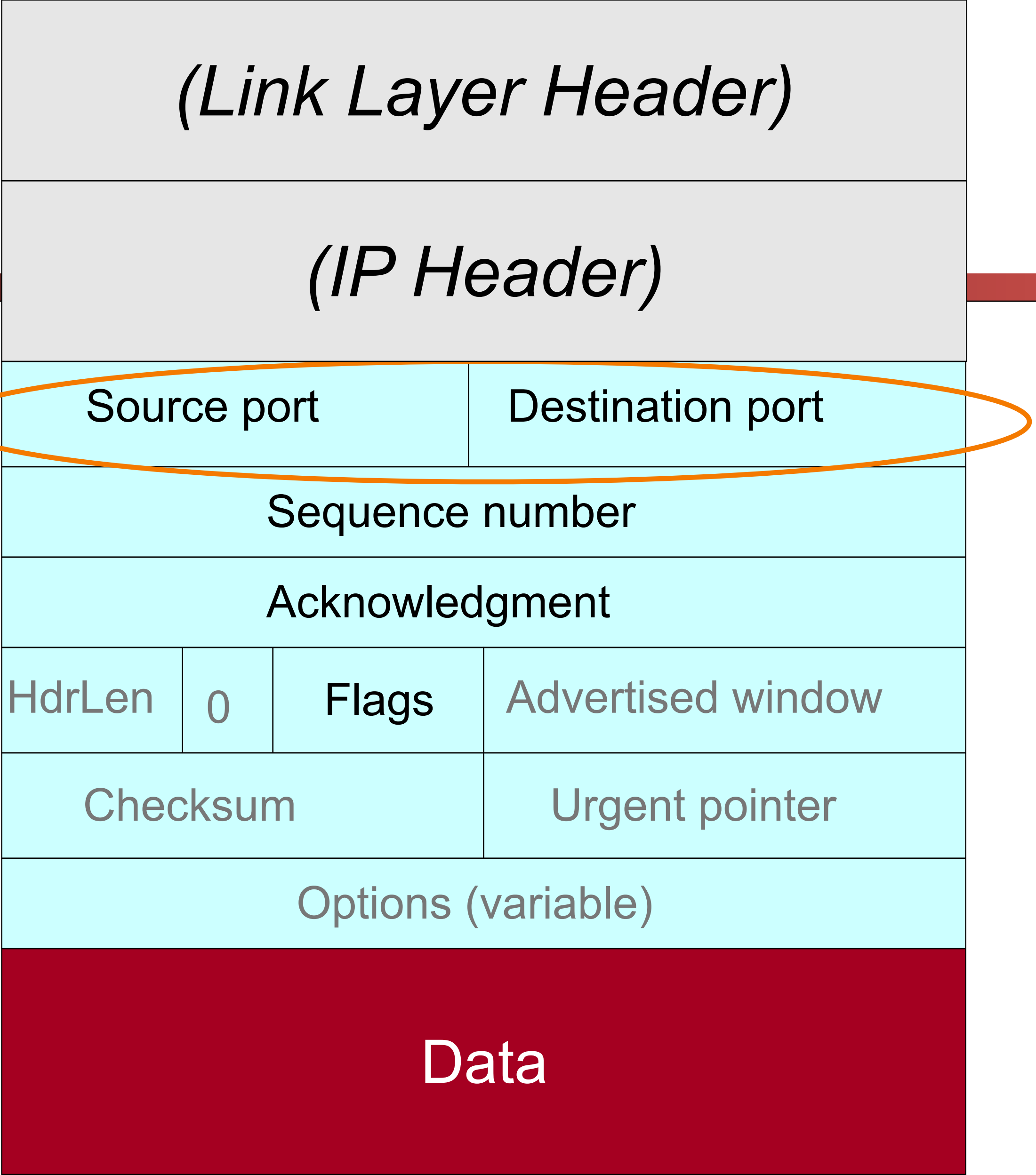
*Ports are associated with OS processes*



# TCP Header

Ports are associated with OS processes

IP source & destination addresses plus TCP source and destination ports uniquely identifies a TCP connection



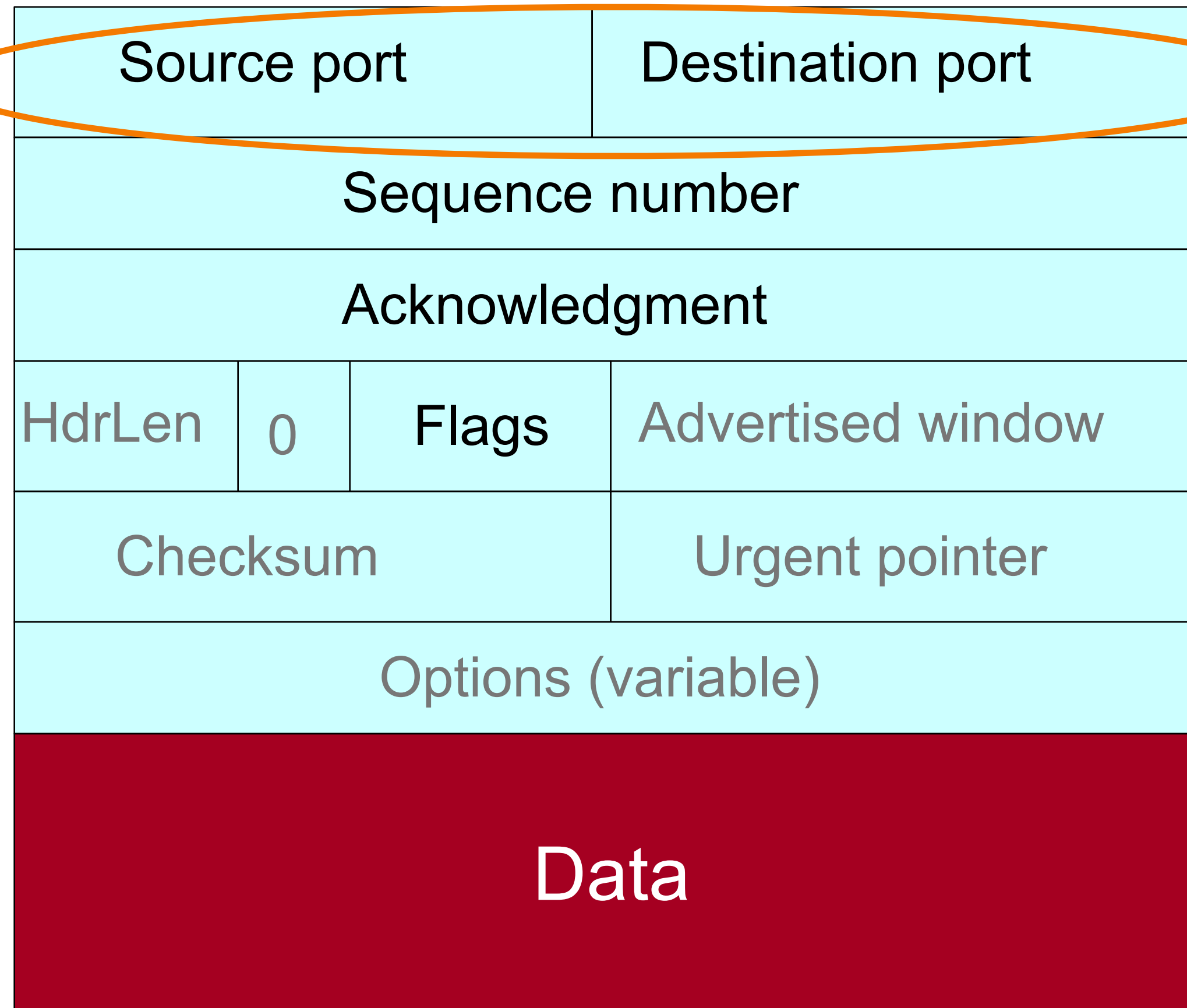


# TCP Header

*Ports are associated with OS processes*

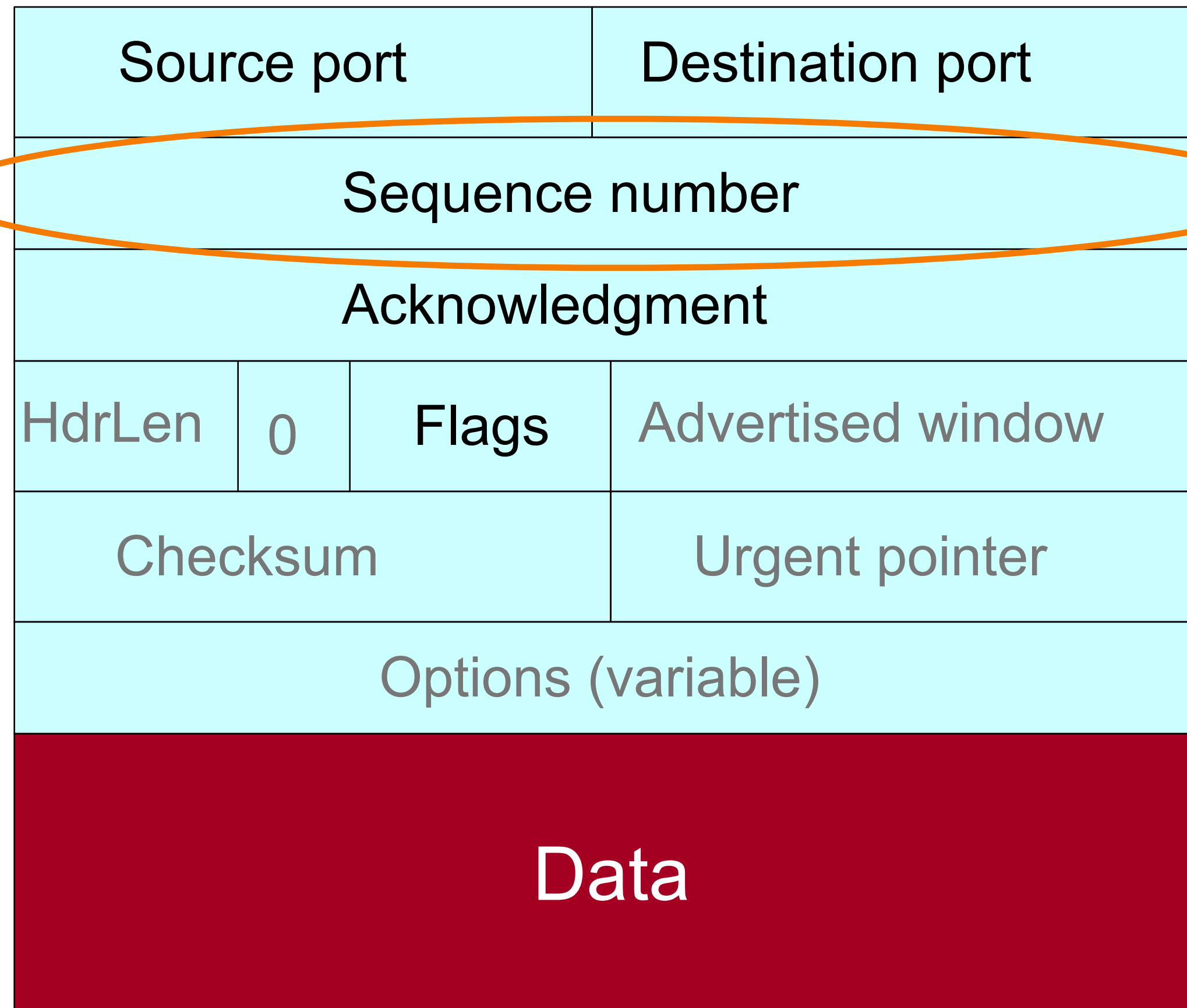
IP source & destination addresses plus TCP source and destination ports uniquely identifies a TCP connection

Some port numbers are “well known” / reserved  
e.g. port 80 = HTTP



# TCP Header

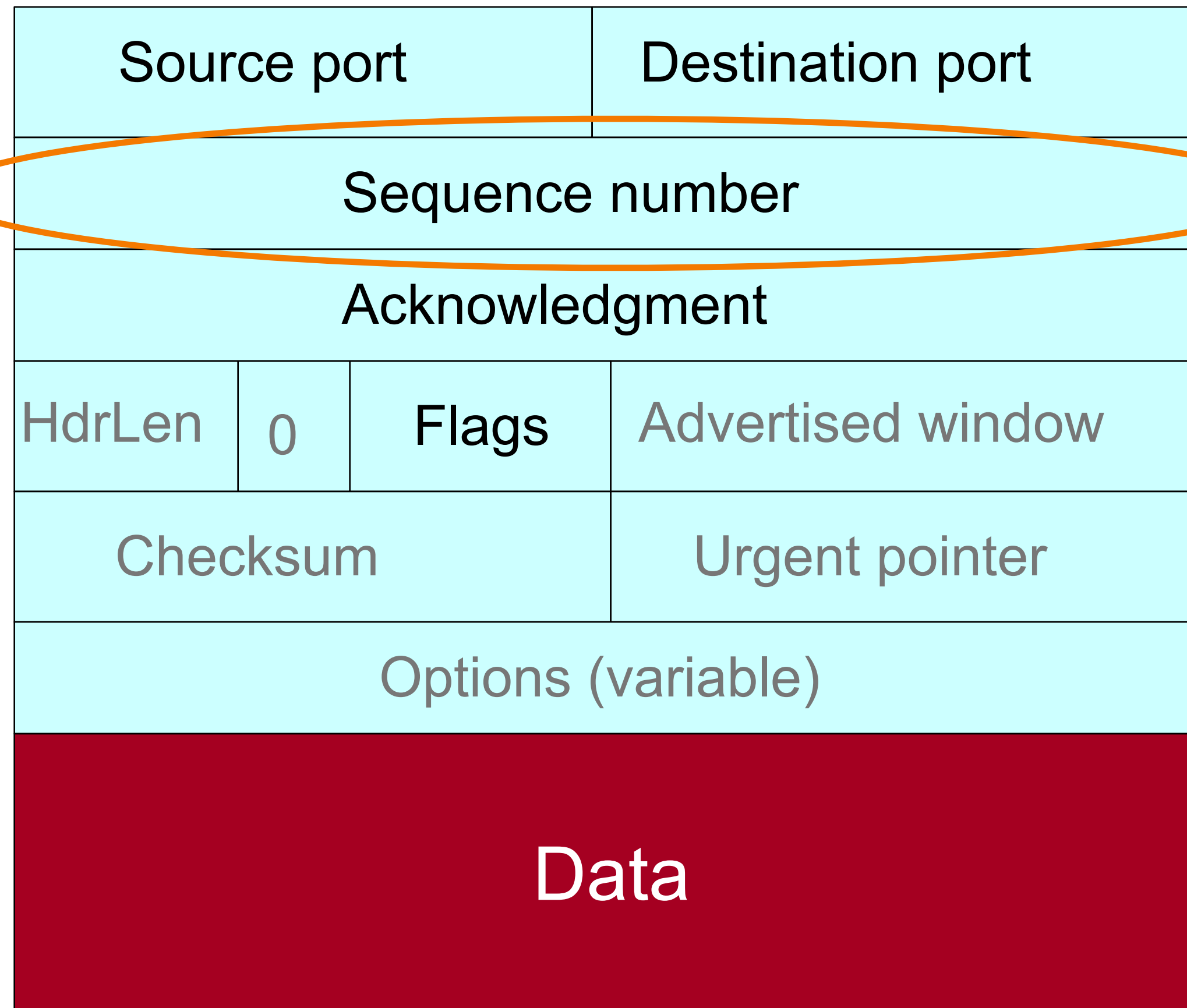
Starting sequence number (byte offset) of data carried in this packet



# TCP Header

Starting sequence number (byte offset) of data carried in this packet

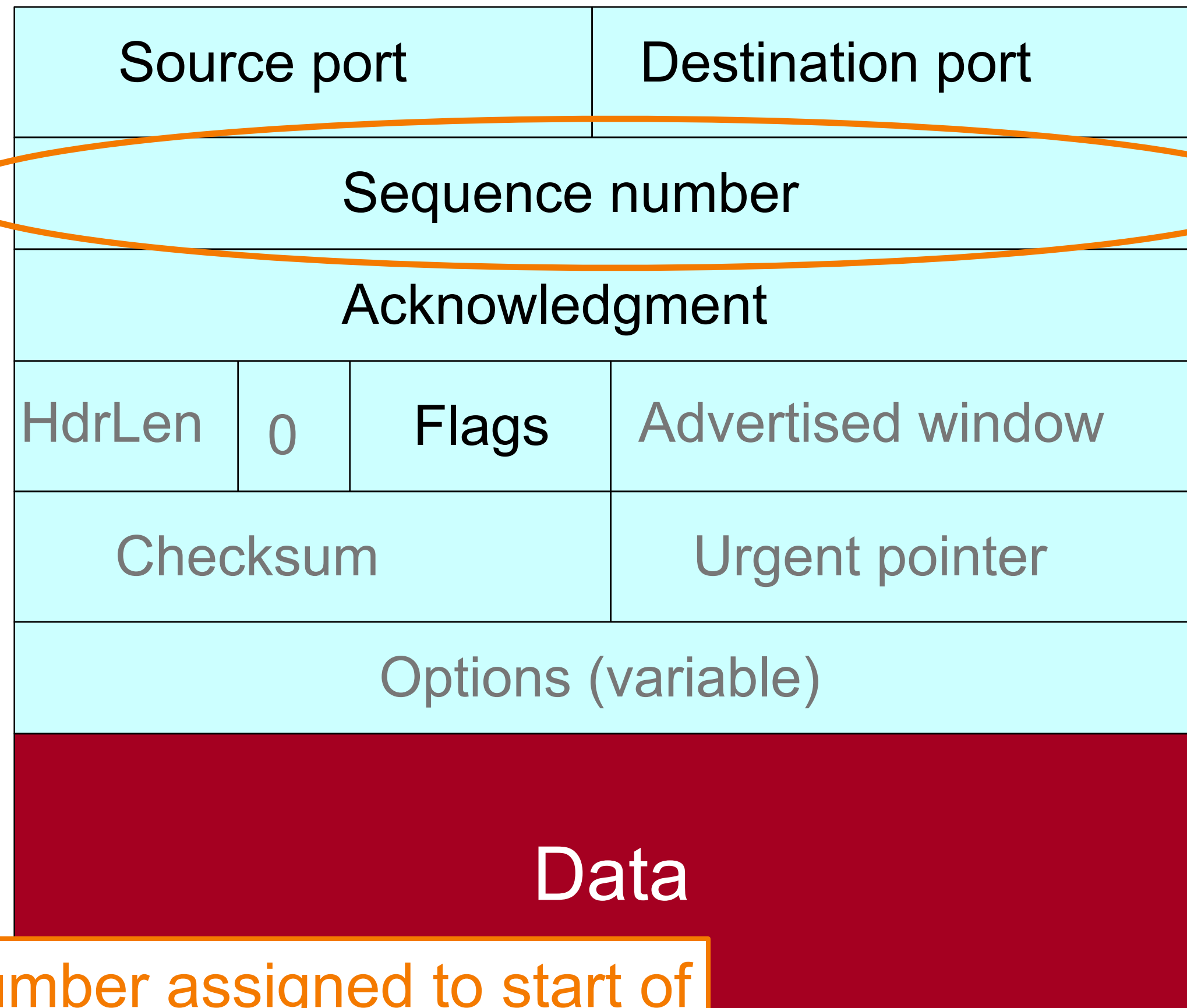
Byte streams numbered independently in each direction



# TCP Header

Starting sequence number (byte offset) of data carried in this packet

Byte stream numbered independently in each direction

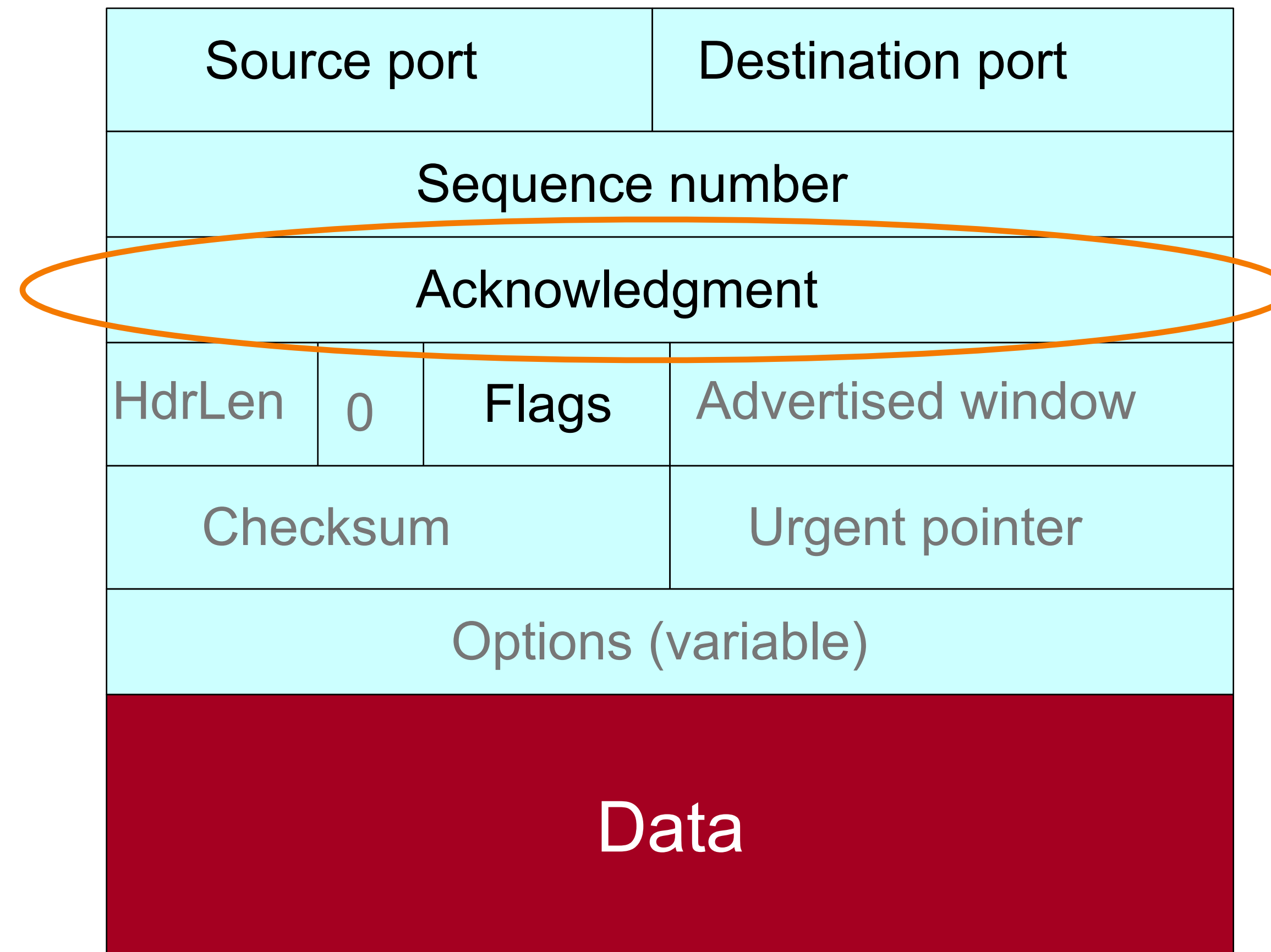


Sequence number assigned to start of byte stream is picked when connection begins; **doesn't** start at 0

# TCP Header

Acknowledgment gives seq # **just beyond** highest seq. received **in order**.

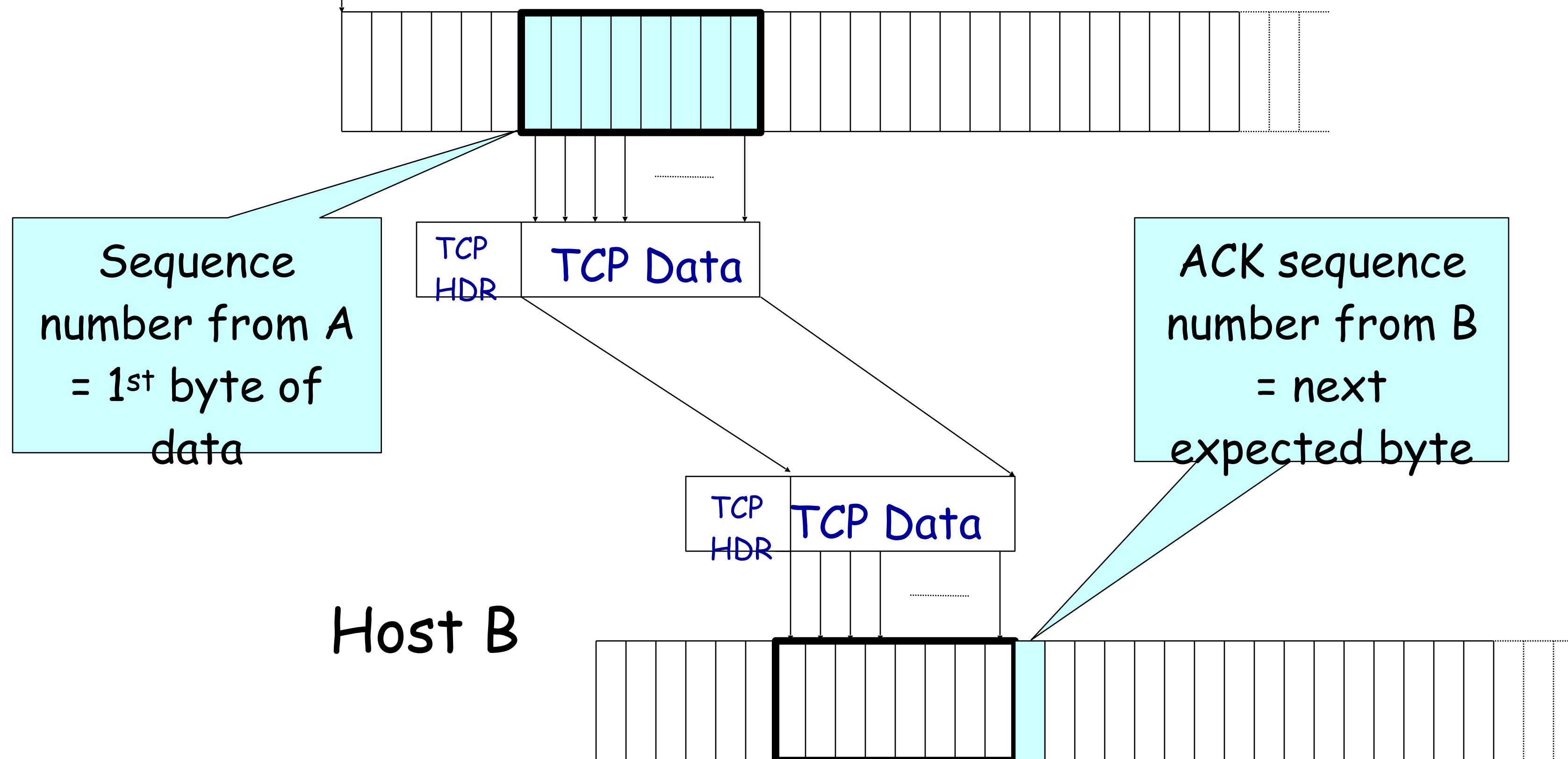
If sender sends **N** bytestream bytes starting at seq **S** then “ack” for it will be **S+N**.



# Sequence Numbers

Host A

ISN (initial sequence number)



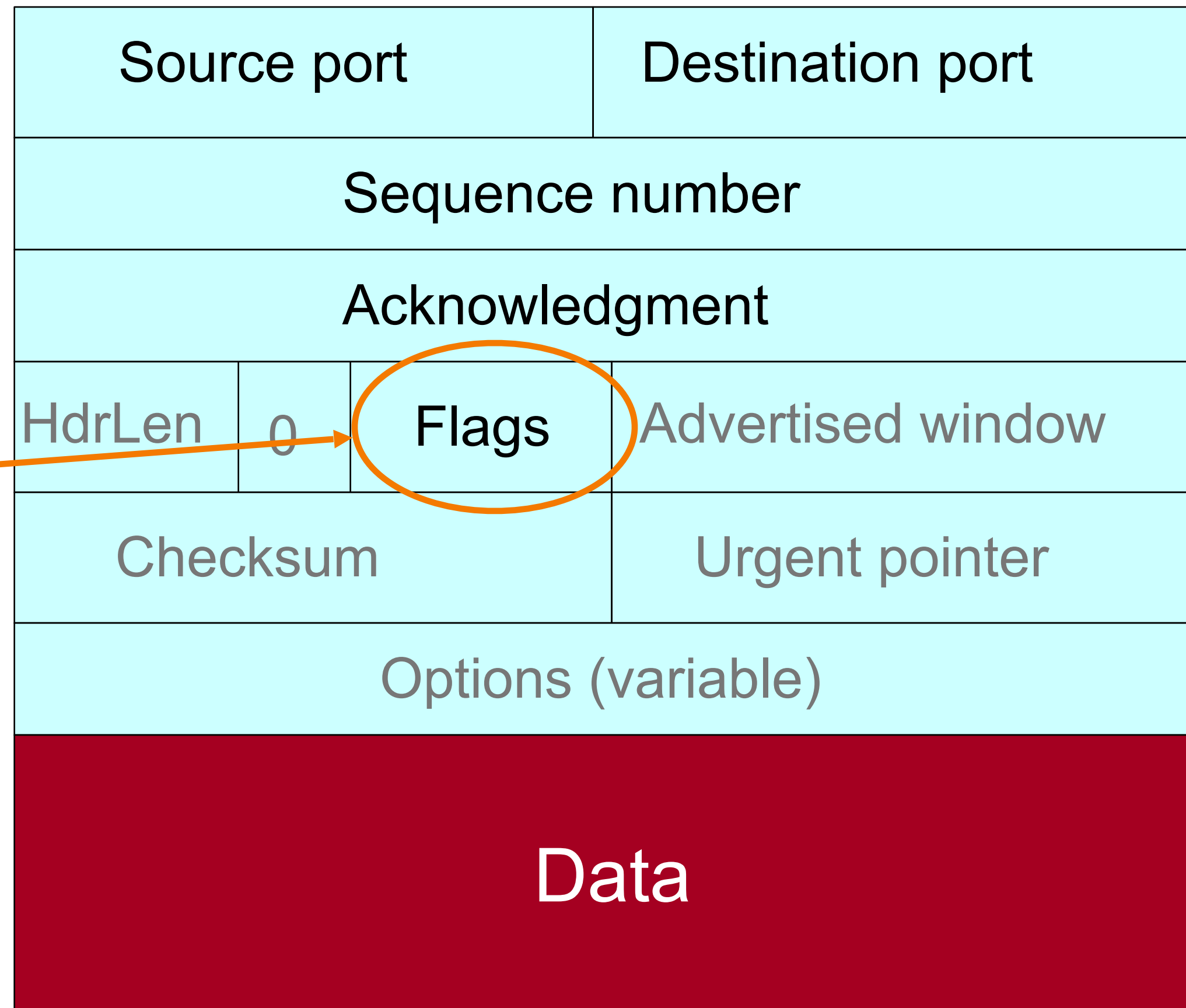
Host B

# TCP Header

Uses include:

acknowledging  
data (“**ACK**”)

setting up (“**SYN**”)  
and closing  
connections  
(“**FIN**” and “**RST**”)

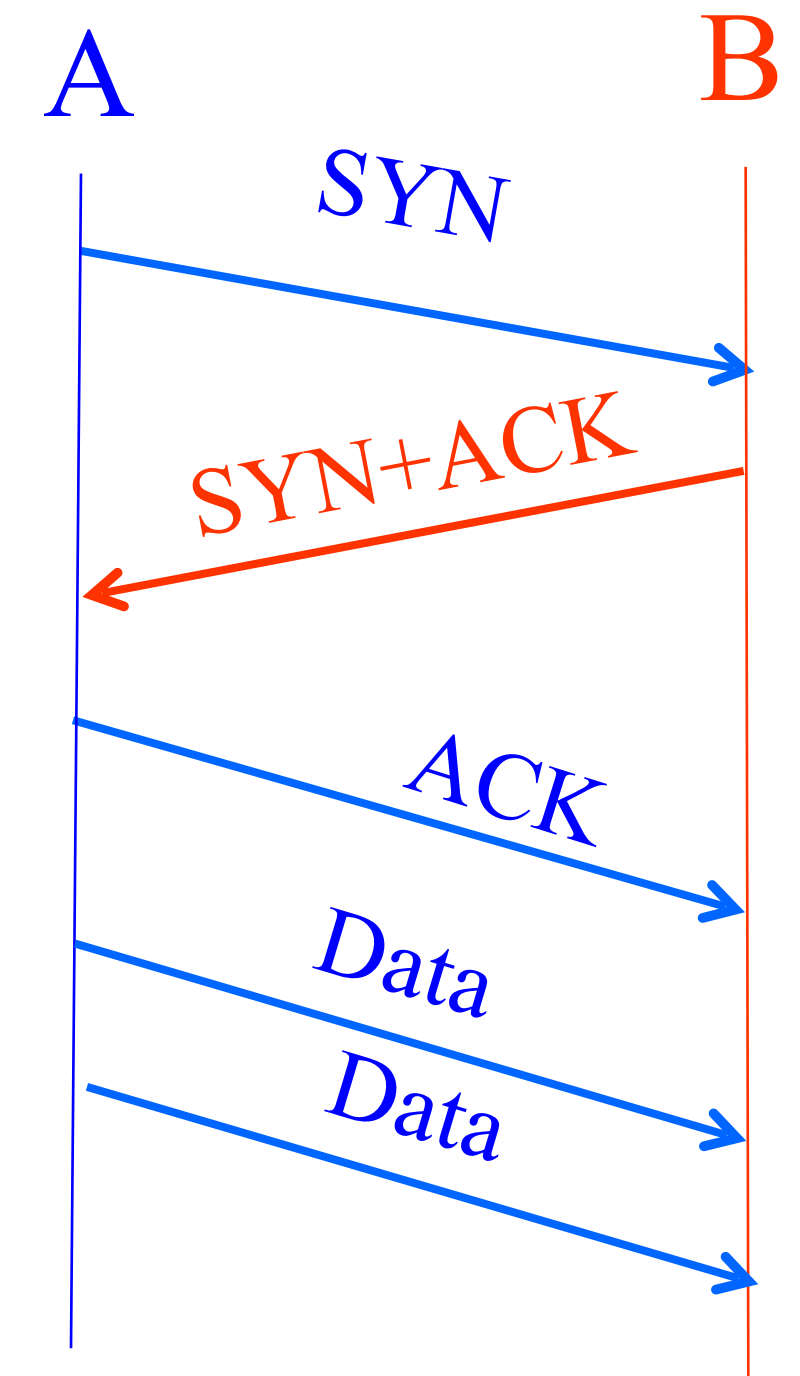


# Establishing a TCP Connection

- Three-way handshake to establish connection
  - Host A sends a **SYN** (open; “synchronize sequence numbers”) to host B
  - Host B returns a SYN acknowledgment (**SYN+ACK**)
  - Host A sends an **ACK** to acknowledge the SYN+ACK

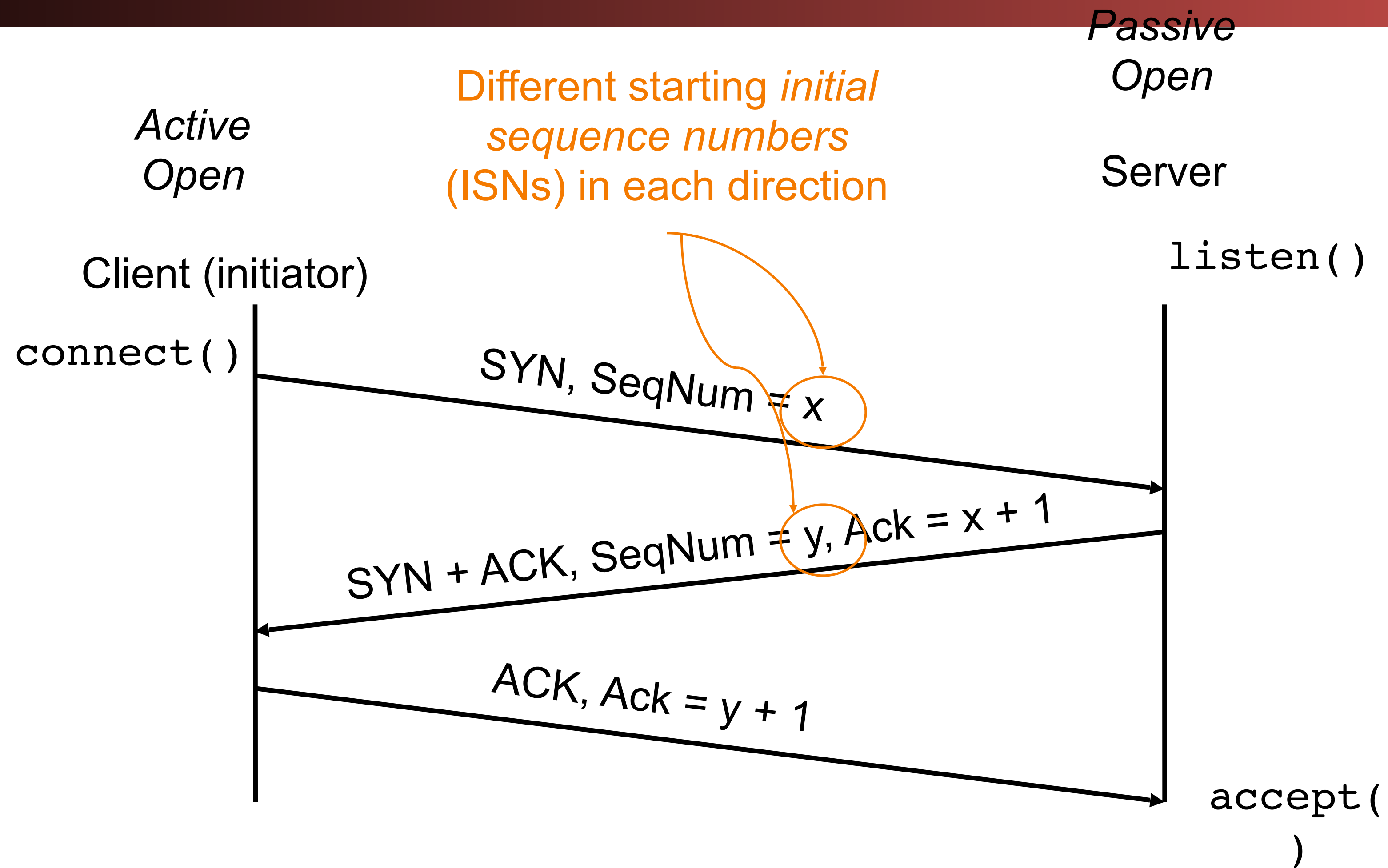
Each host tells its *Initial Sequence Number* (ISN) to the other host.

(Spec says to pick based on local clock)





# Timing Diagram: 3-Way Handshaking



# Host Names vs. IP addresses

- Host names
  - Examples: `www.cnn.com` and `bbc.co.uk`
  - Mnemonic name appreciated by **humans**
  - Variable length, full alphabet of characters
  - Provide little (if any) information about location
- IP addresses
  - Examples: `64.236.16.20` and `212.58.224.131`
  - Numerical address appreciated by **routers**
  - Fixed length, binary number
  - Hierarchical, related to host location

# So Let's Do A Google Search...

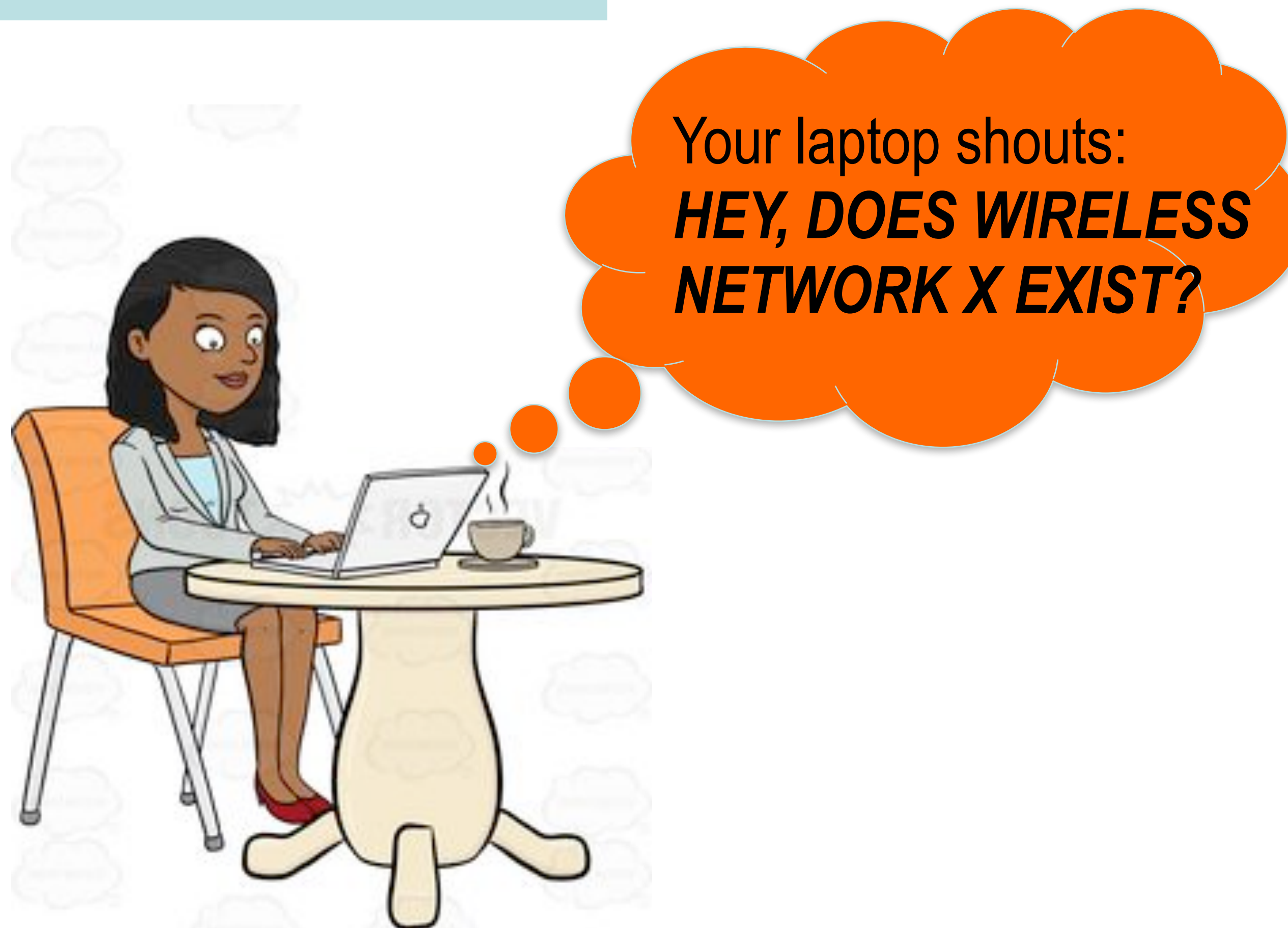
- Walk into a coffee shop
- Open a laptop
- Search google...

# Coffee Shop



# Coffee Shop

## 1. Join the wireless network





# Coffee Shop

## 1. Join the wireless network

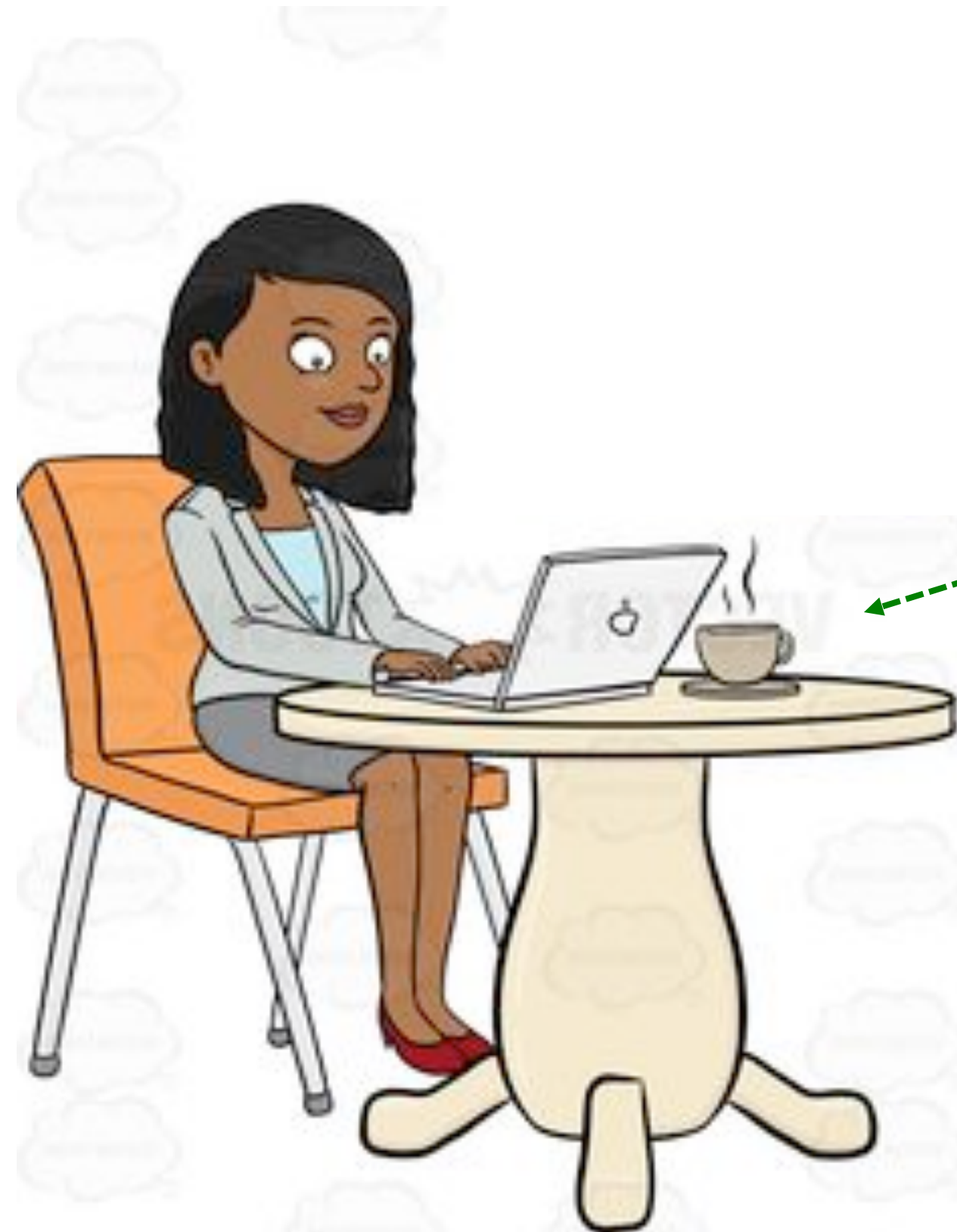


Wireless access point(s)  
continually shout:  
***HEY, I'M WIRELESS  
NETWORK Y, JOIN ME!***



# Coffee Shop

## 1. Join the wireless network



If either match up, your laptop joins the network. Optionally performs a cryptographic exchange.





# Coffee Shop

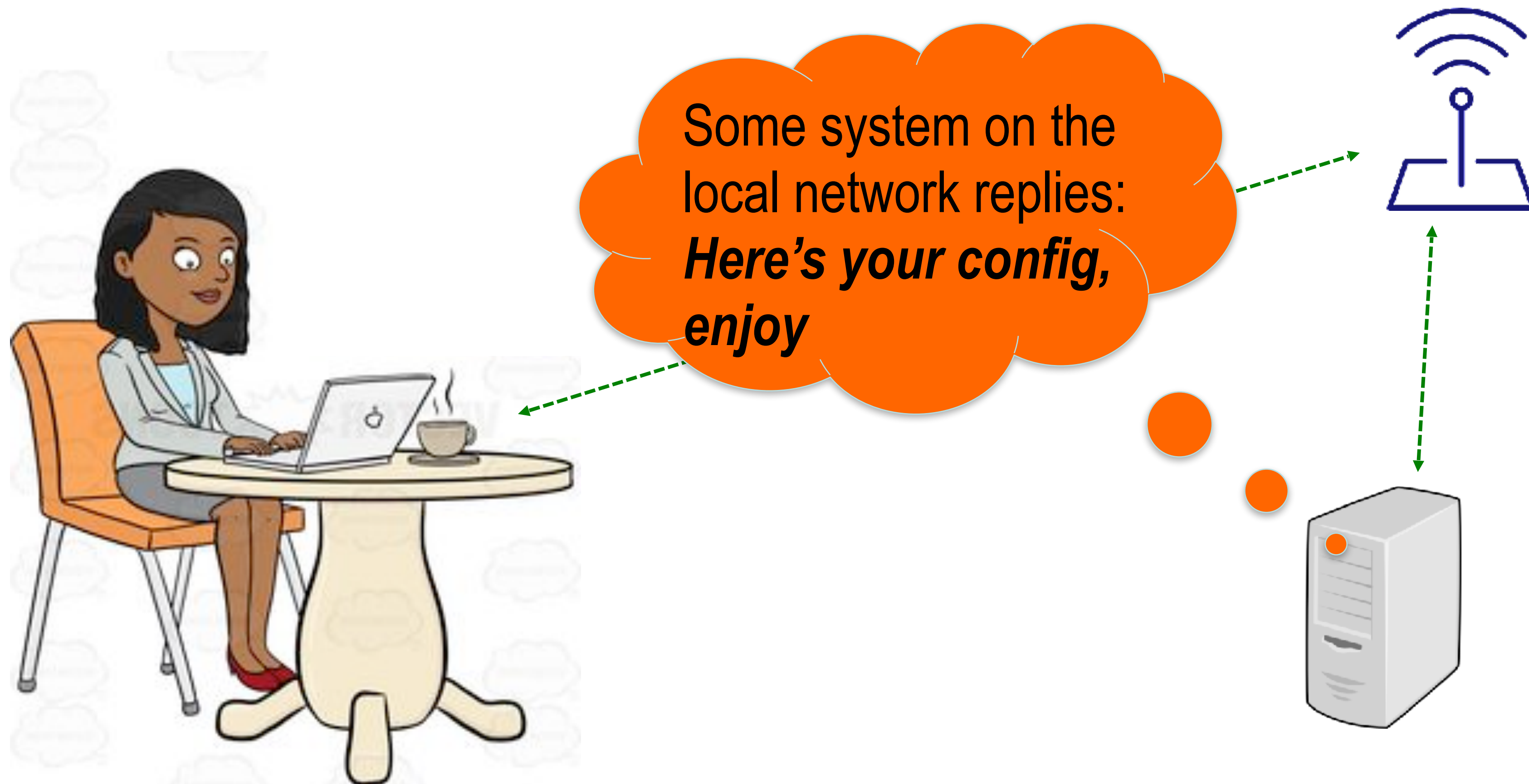
## 2. Configure your connection





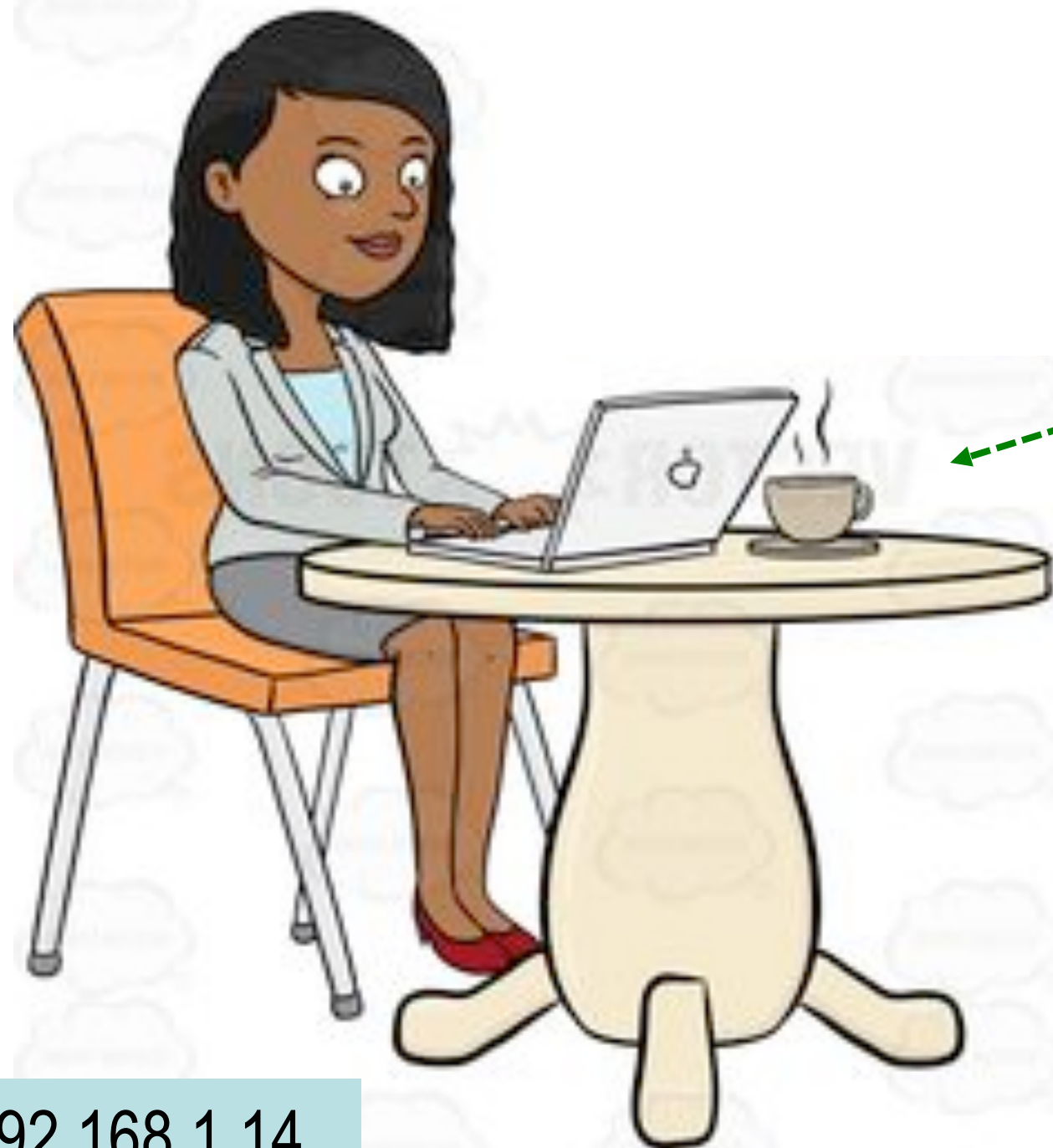
# Coffee Shop

## 2. Configure your connection



# Coffee Shop

## 2. Configure your connection



192.168.1.14

The configuration includes:

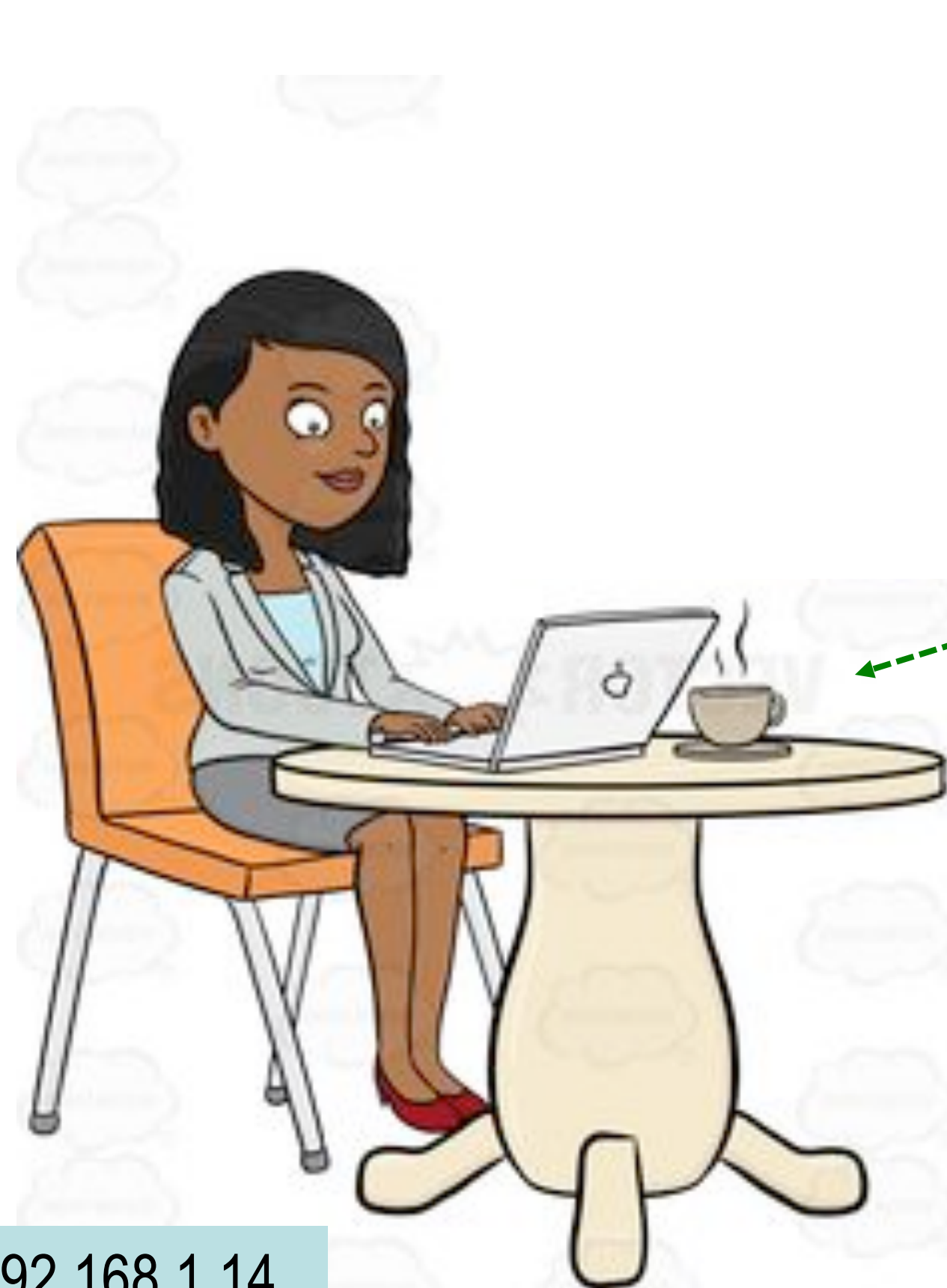
- (1) An Internet address (**IP address**) your laptop should use; typically 32 bits (IPv4).
- (2) The address of a “**gateway**” system to use to access *hosts* beyond the local network
- (3) The address of a **DNS server** (“*resolver*”) to map names like `google.com` to IP addresses





# Coffee Shop

## 3. Find the address of `google.com`



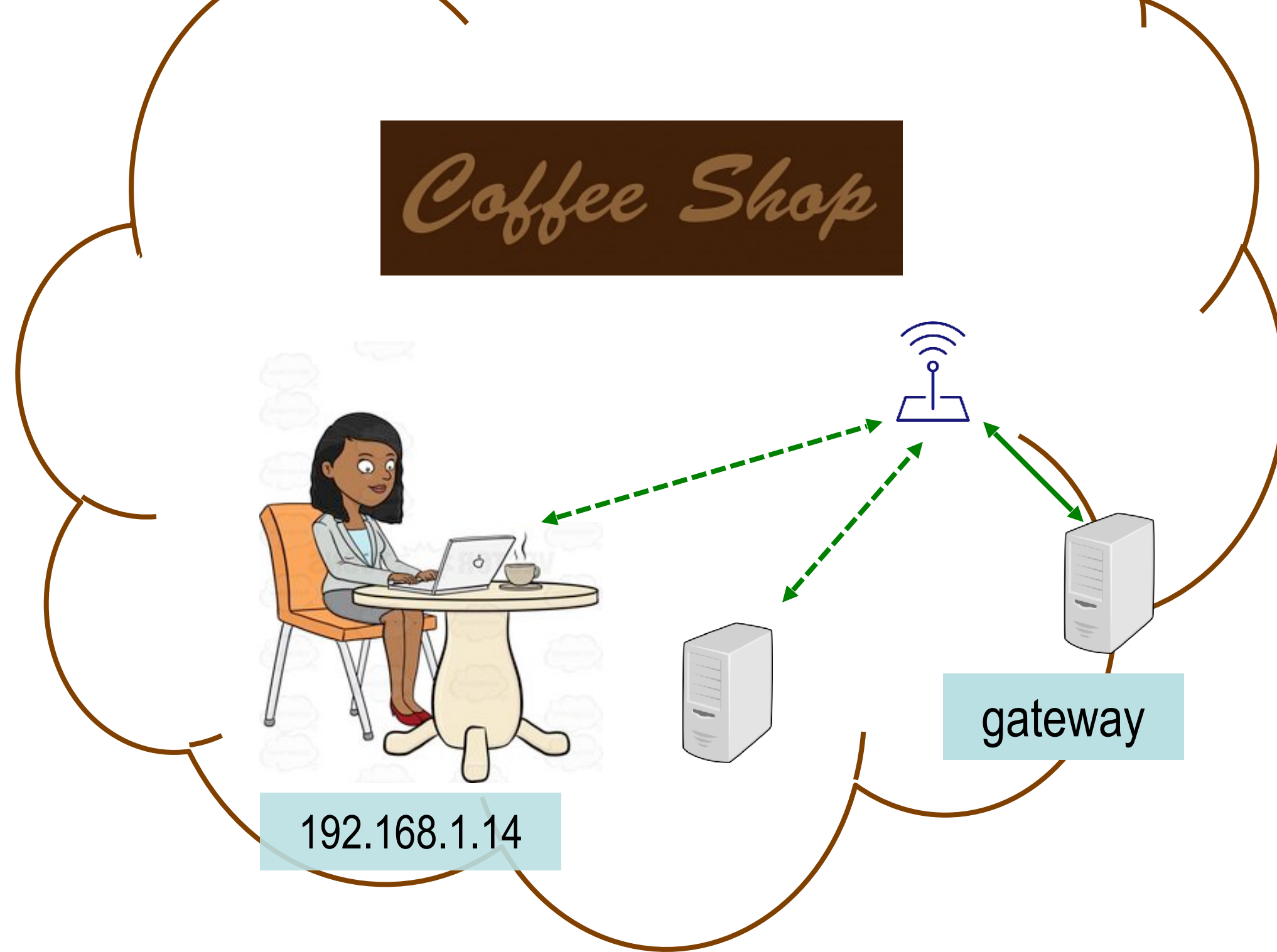
Your laptop sends a **DNS** request asking: “*address for google.com?*”

It's transmitted using the **UDP** protocol (lightweight, unreliable).

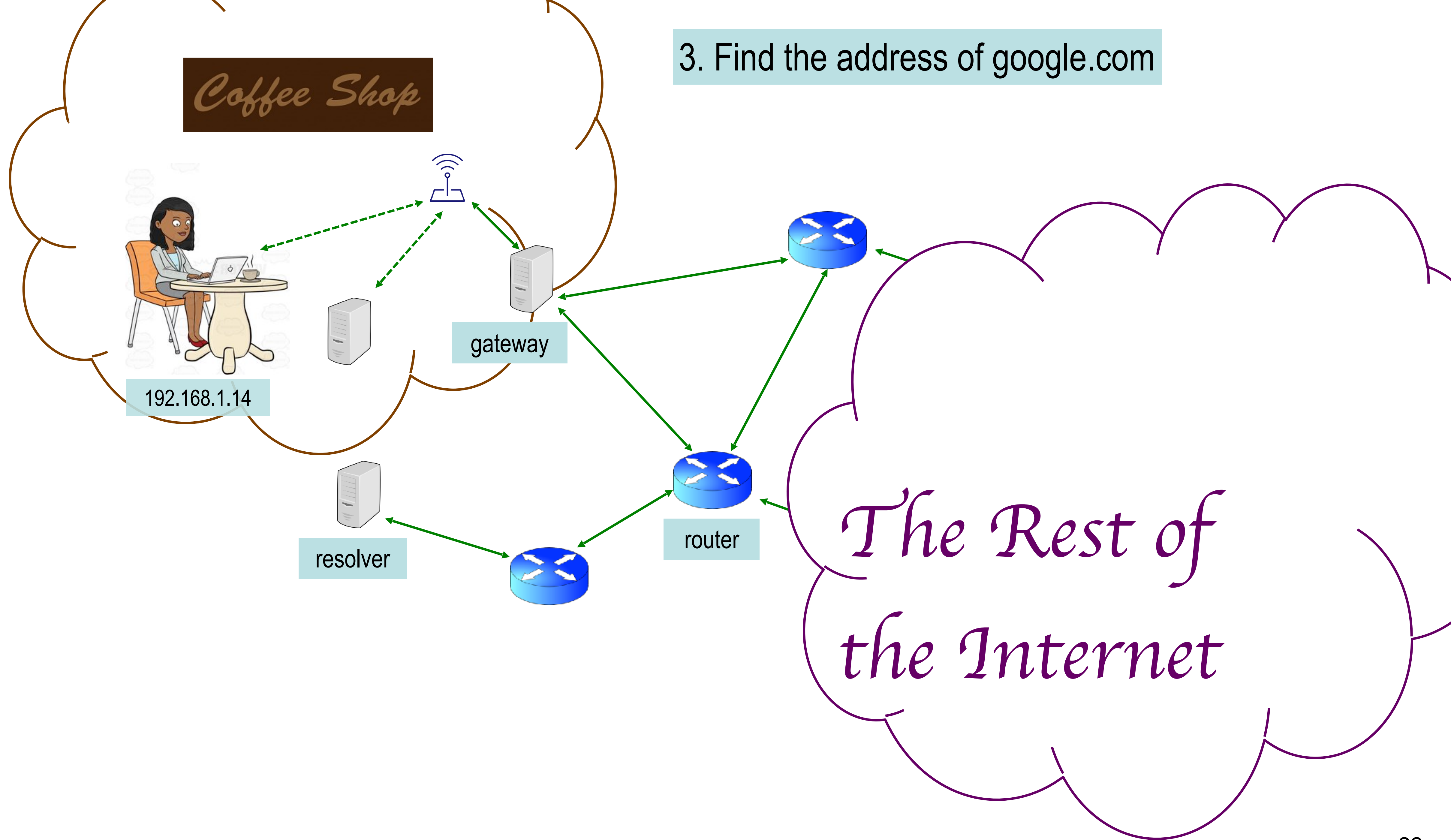
The DNS **resolver** might not be on the local network.

192.168.1.14

### 3. Find the address of google.com

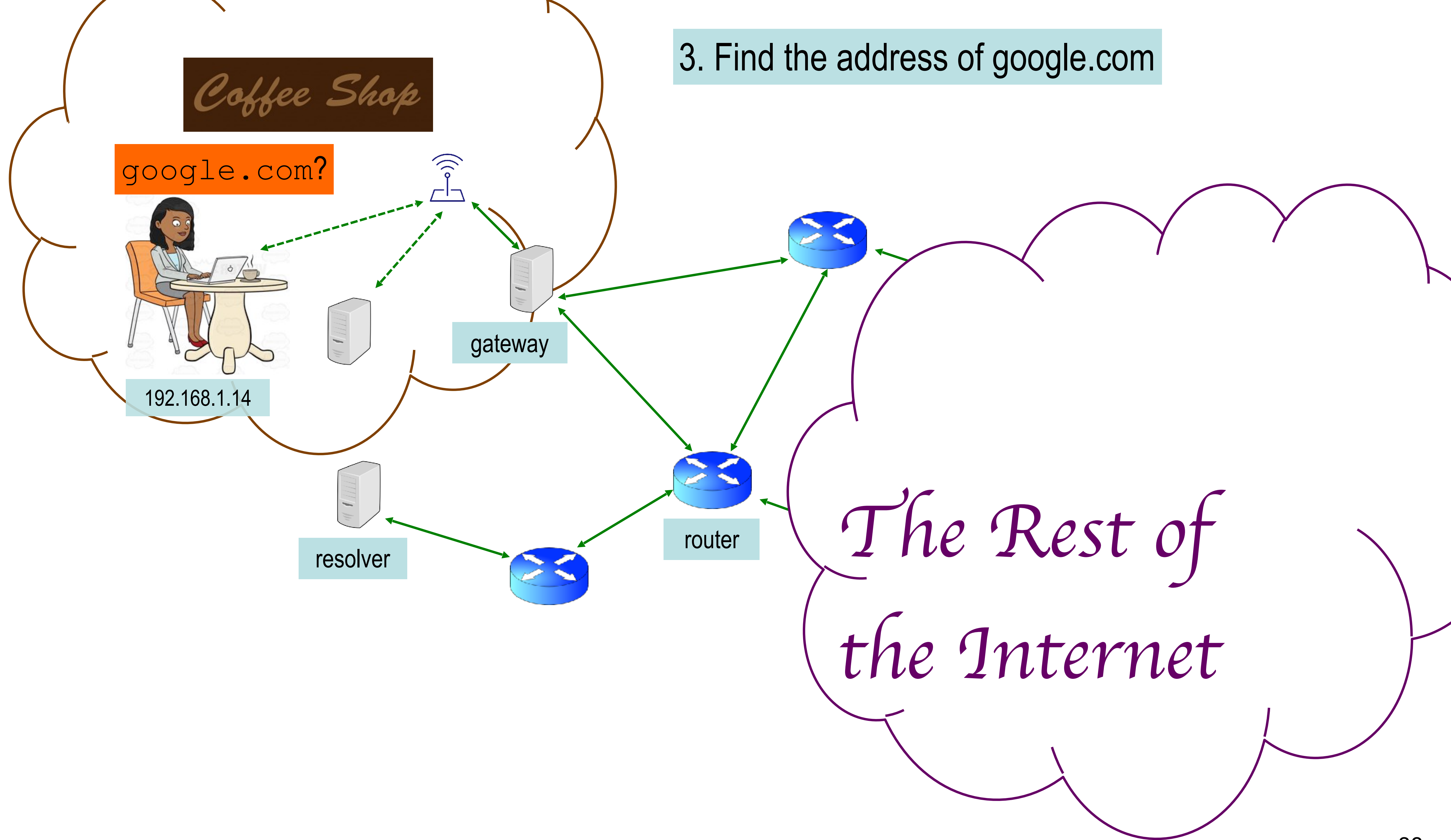


3. Find the address of google.com

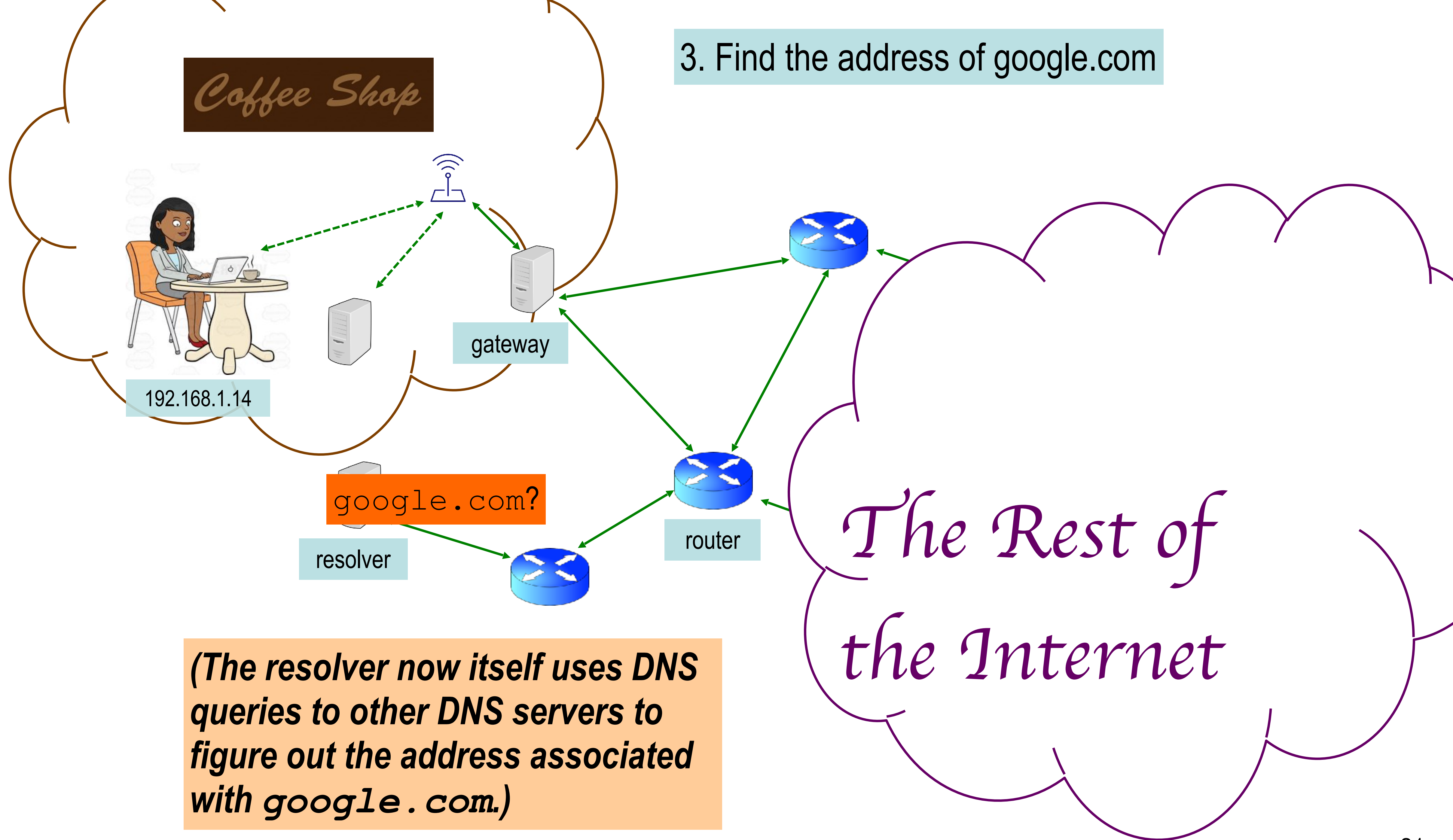




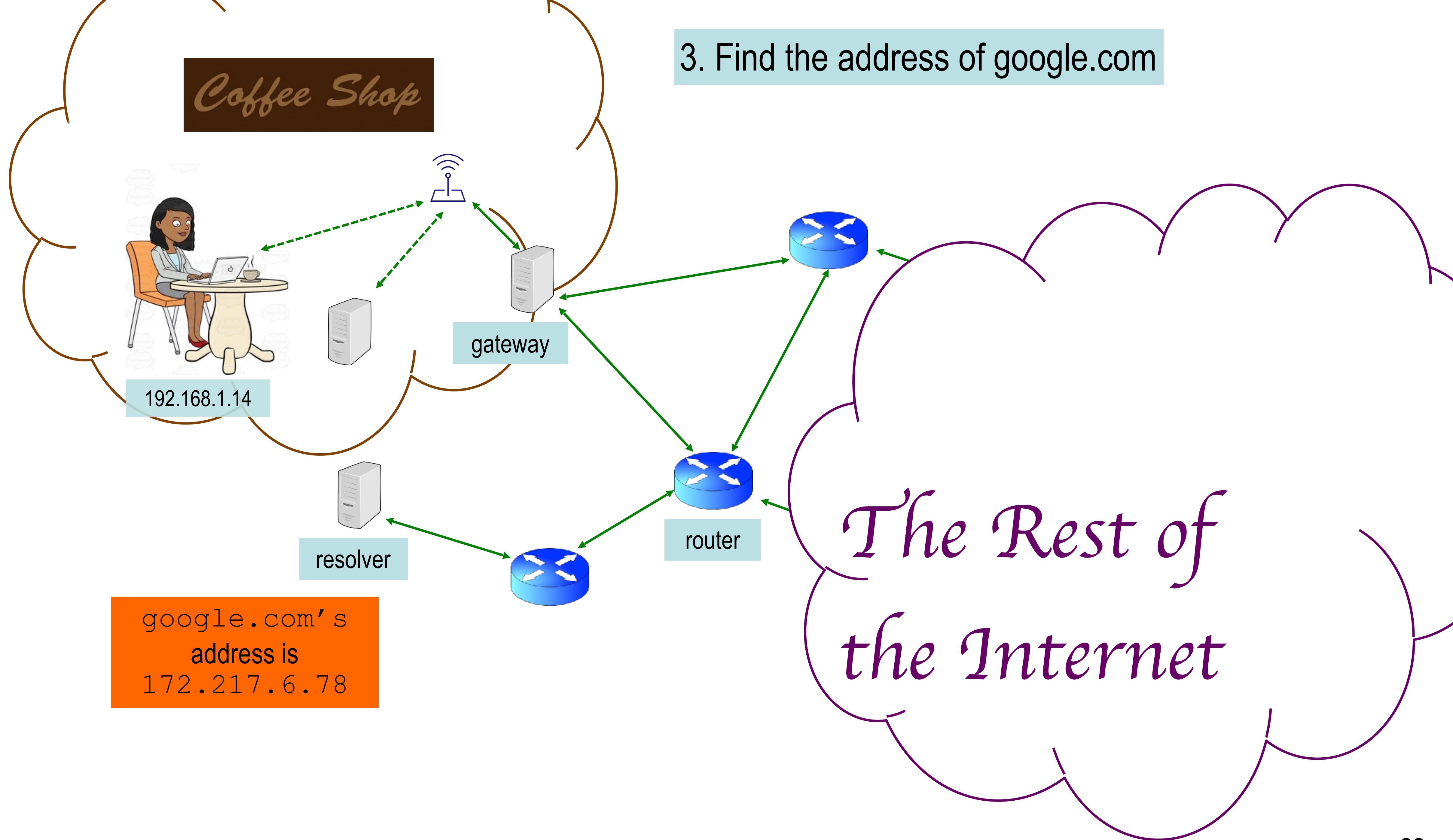
3. Find the address of google.com



3. Find the address of google.com



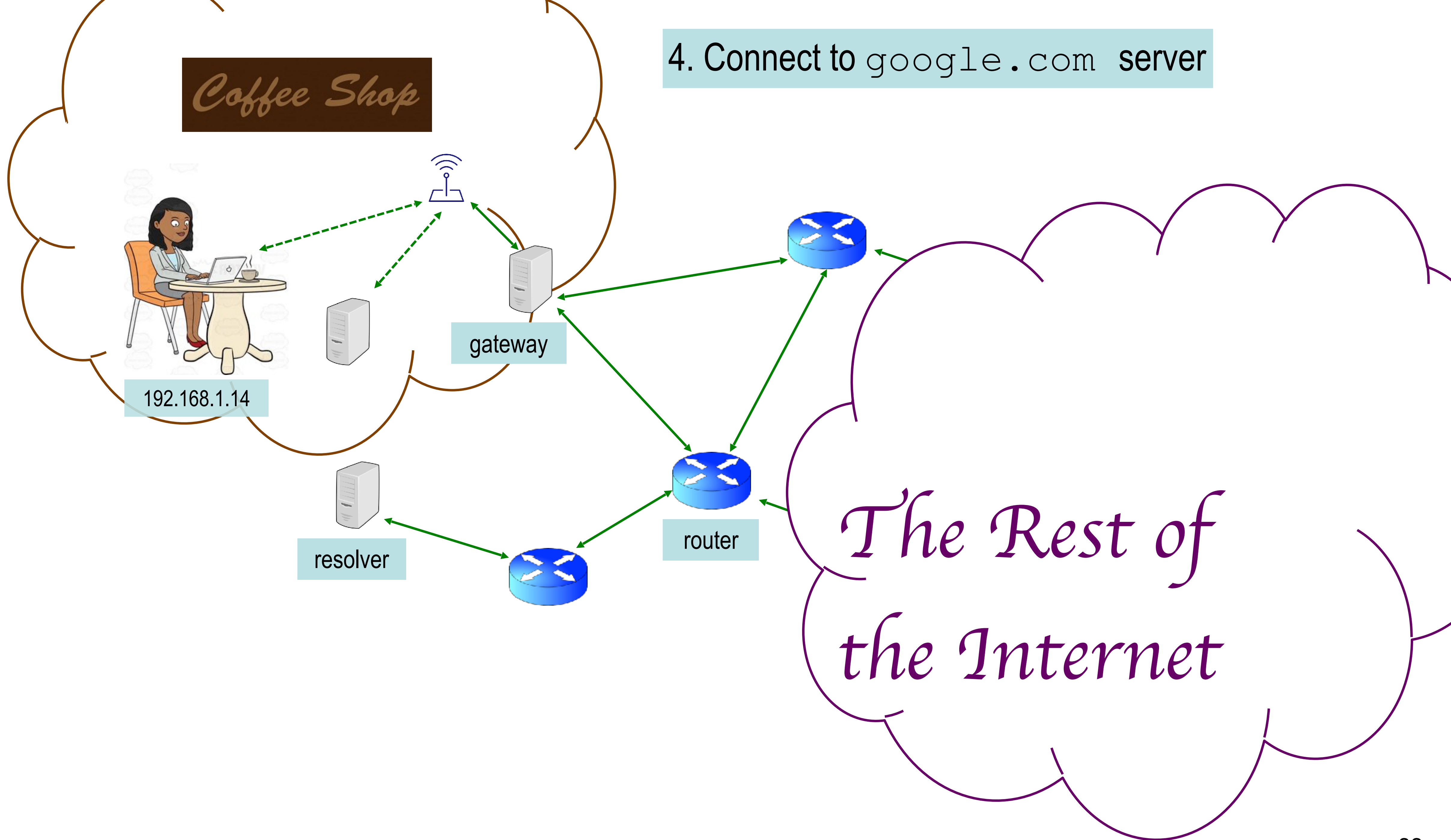
3. Find the address of google.com



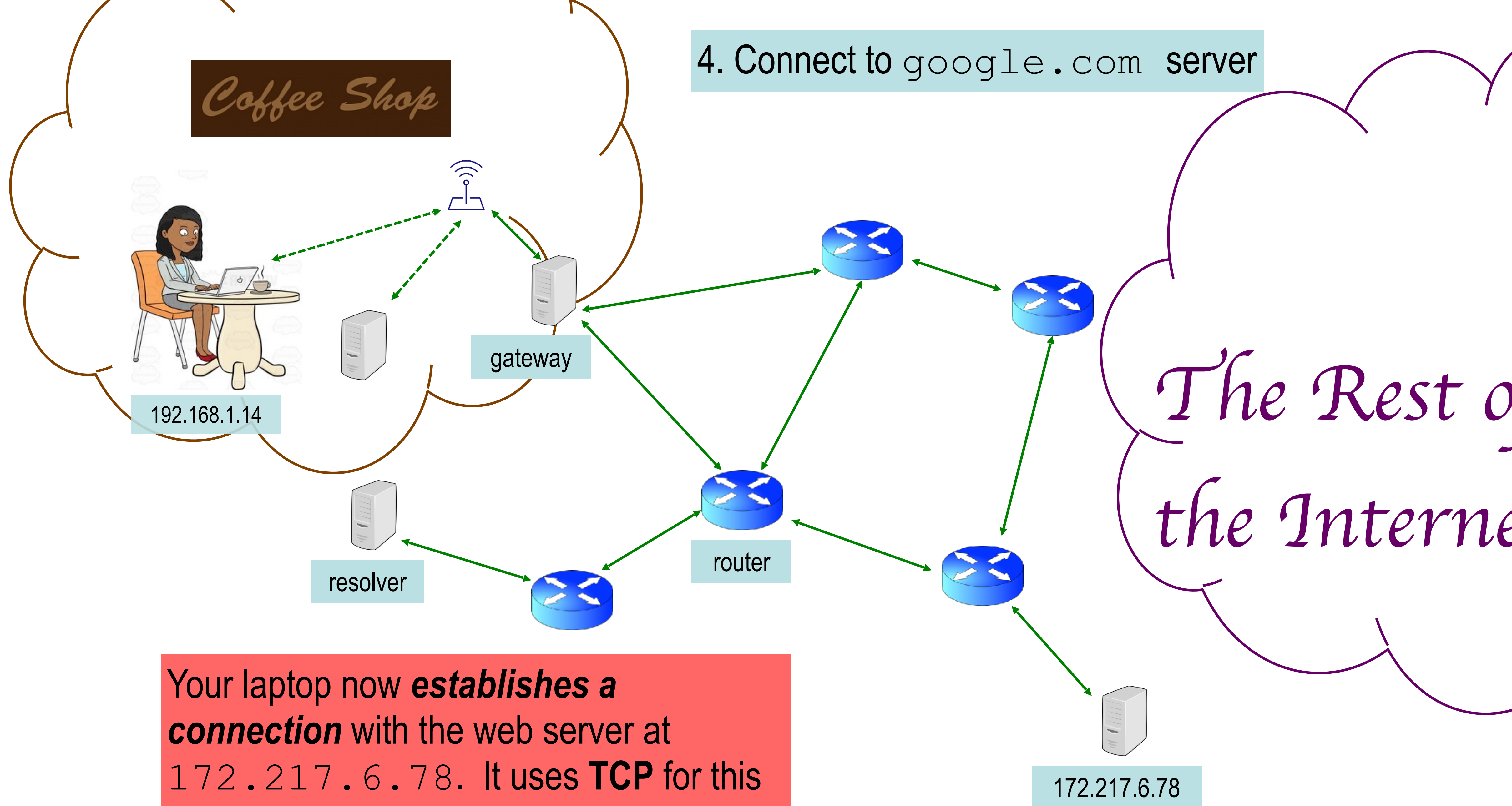
google.com's  
address is  
172.217.6.78



#### 4. Connect to google.com server

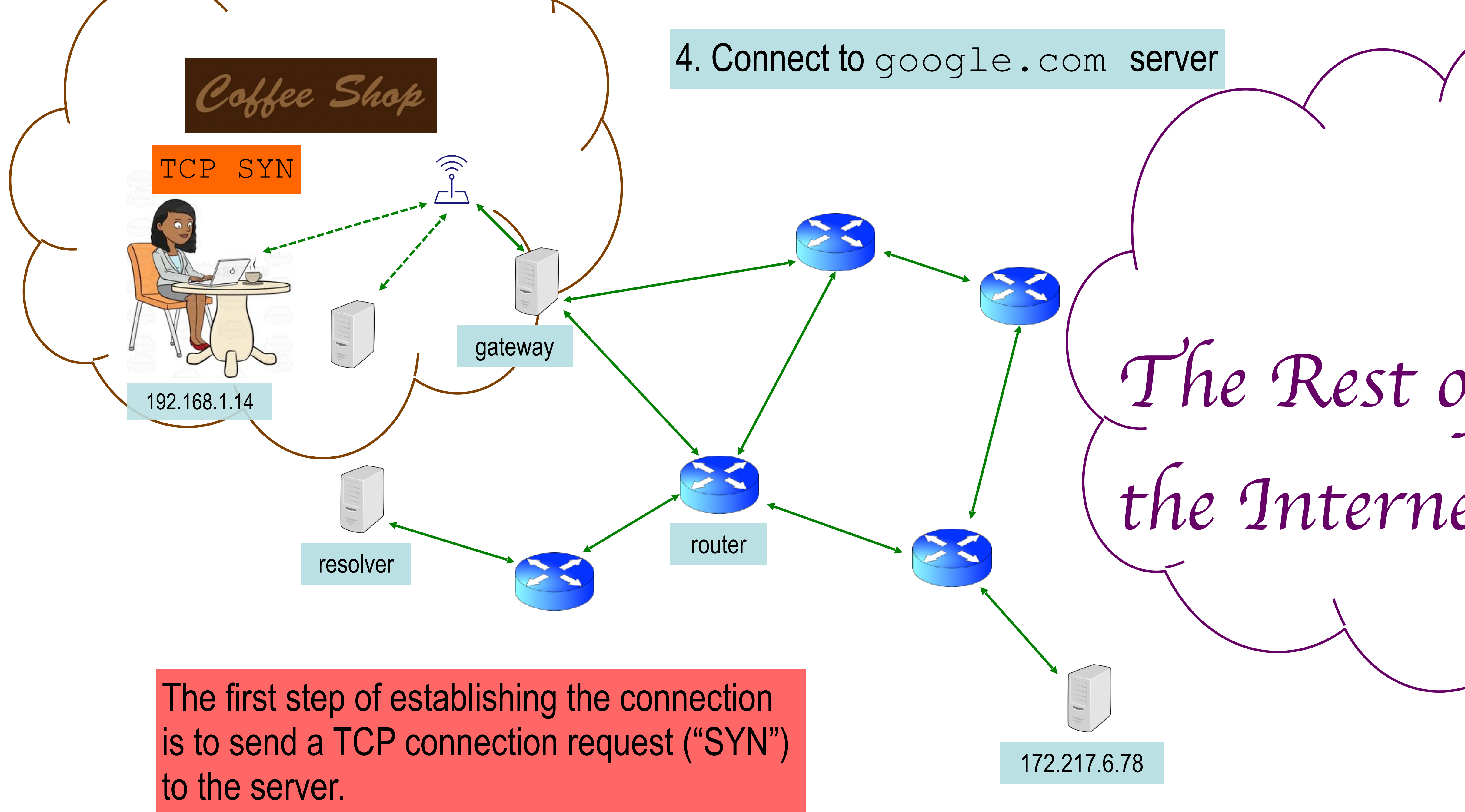


#### 4. Connect to google.com server



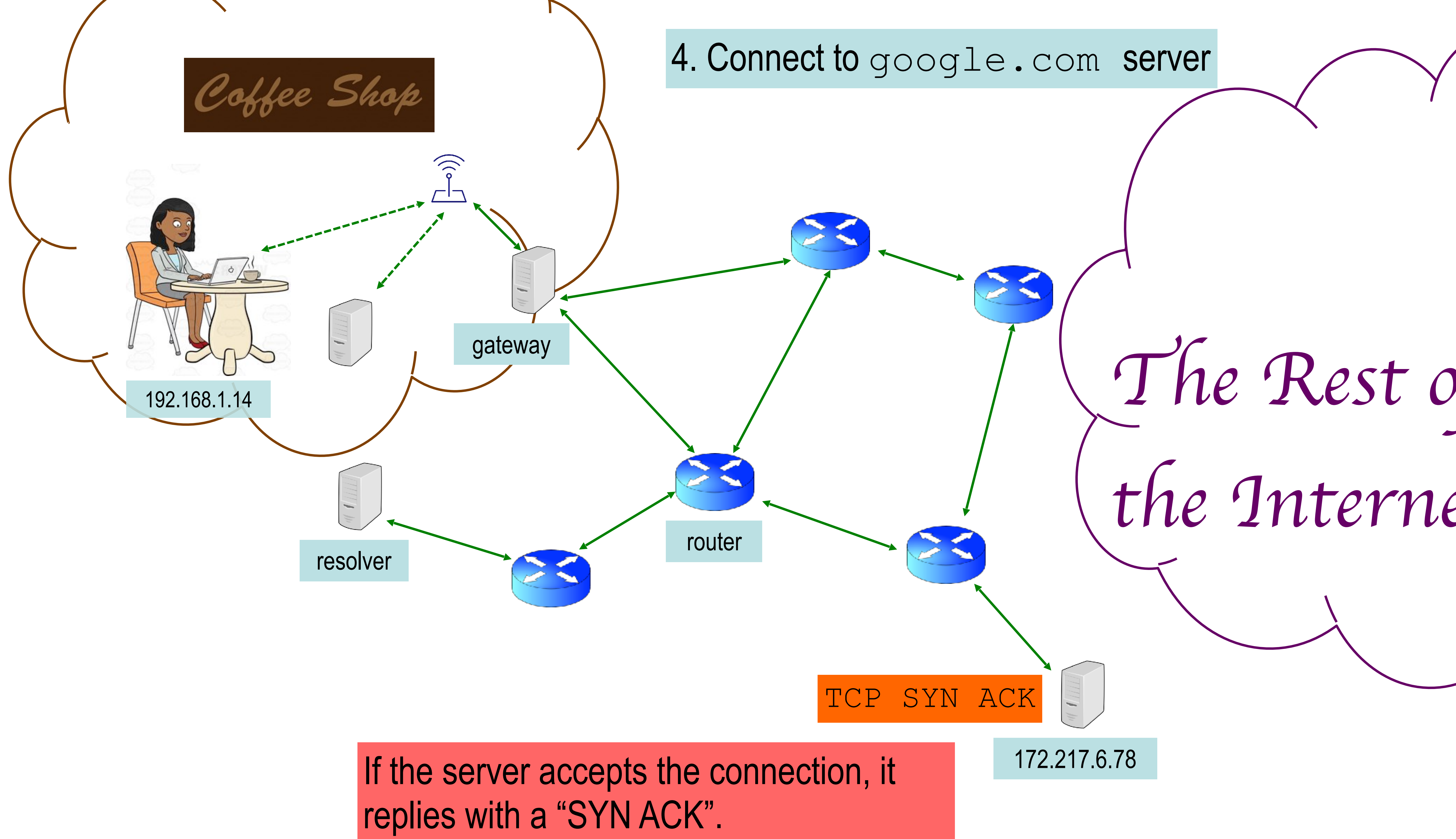
Your laptop now **establishes a connection** with the web server at 172.217.6.78. It uses **TCP** for this rather than UDP, to obtain reliability.

#### 4. Connect to google.com server

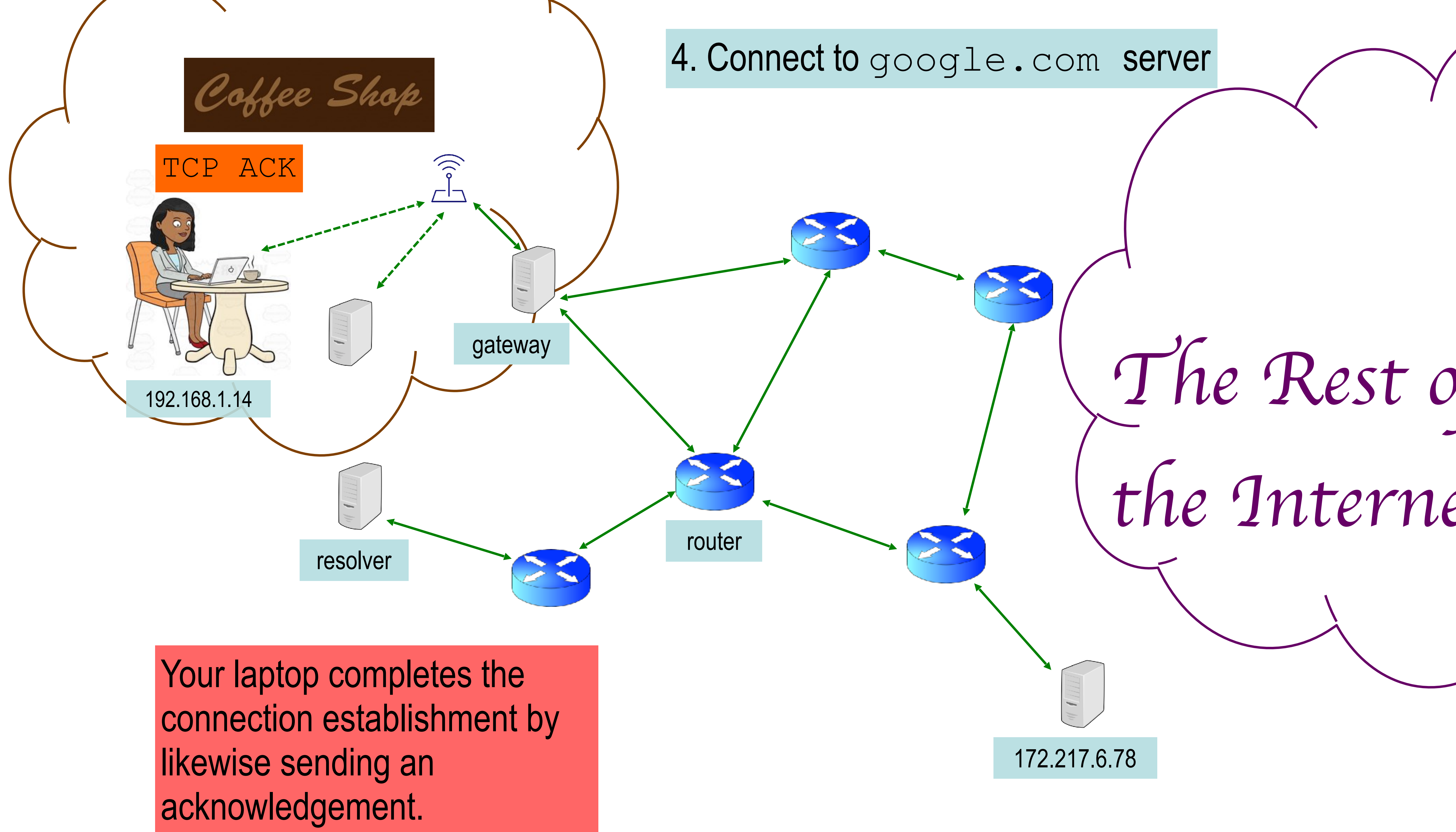




#### 4. Connect to google.com server

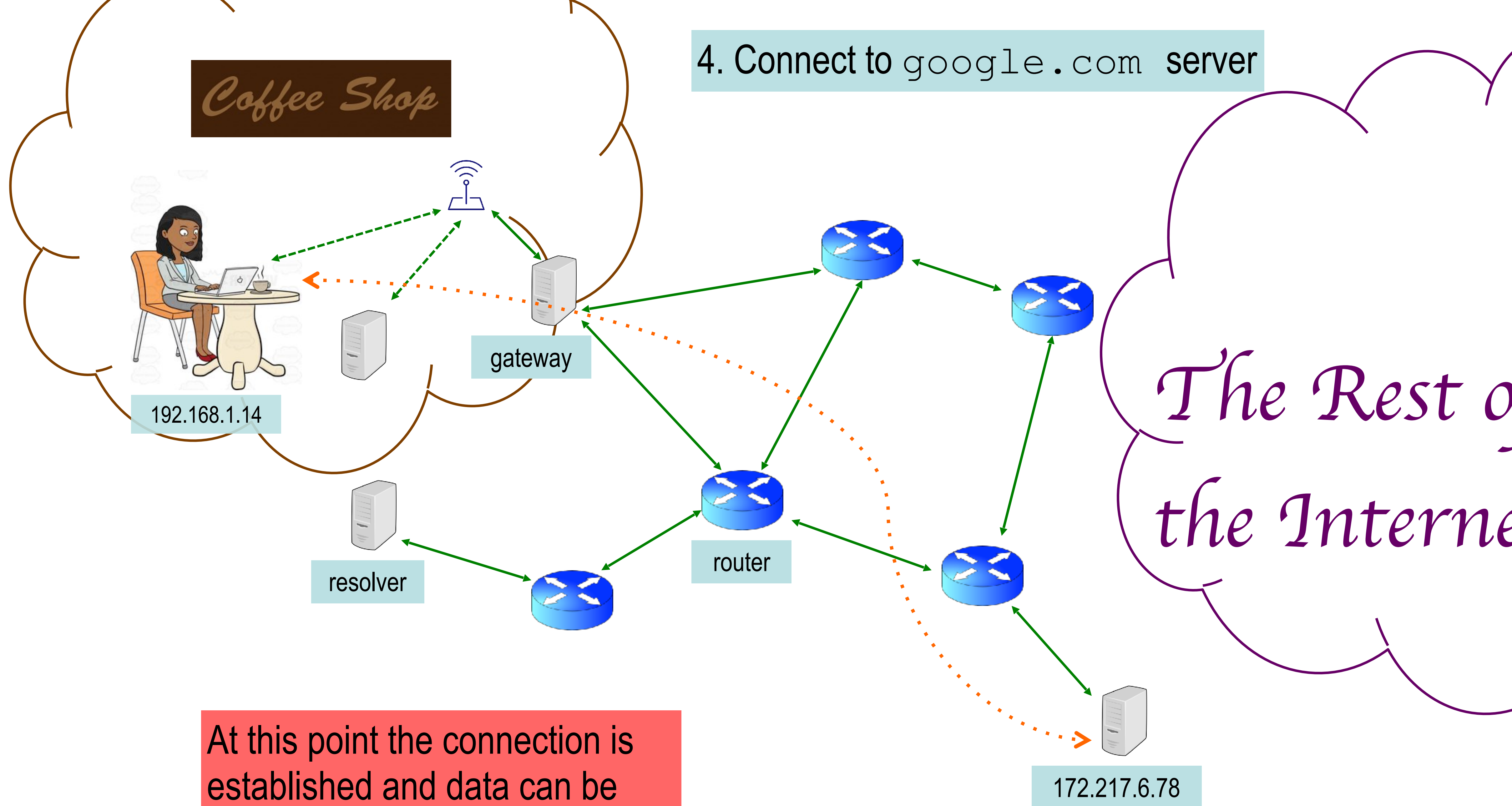


#### 4. Connect to google.com server



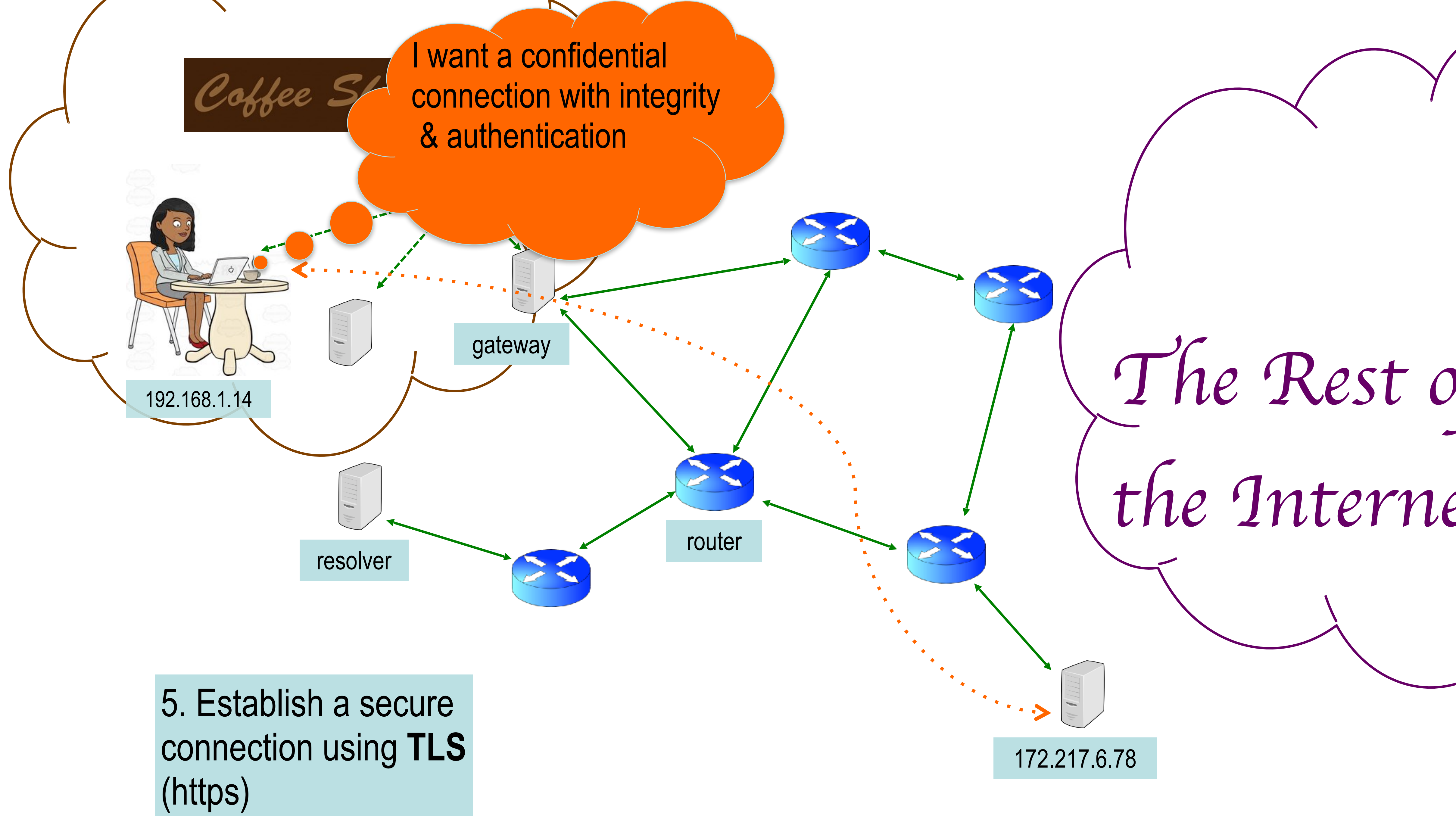
Your laptop completes the connection establishment by likewise sending an acknowledgement.

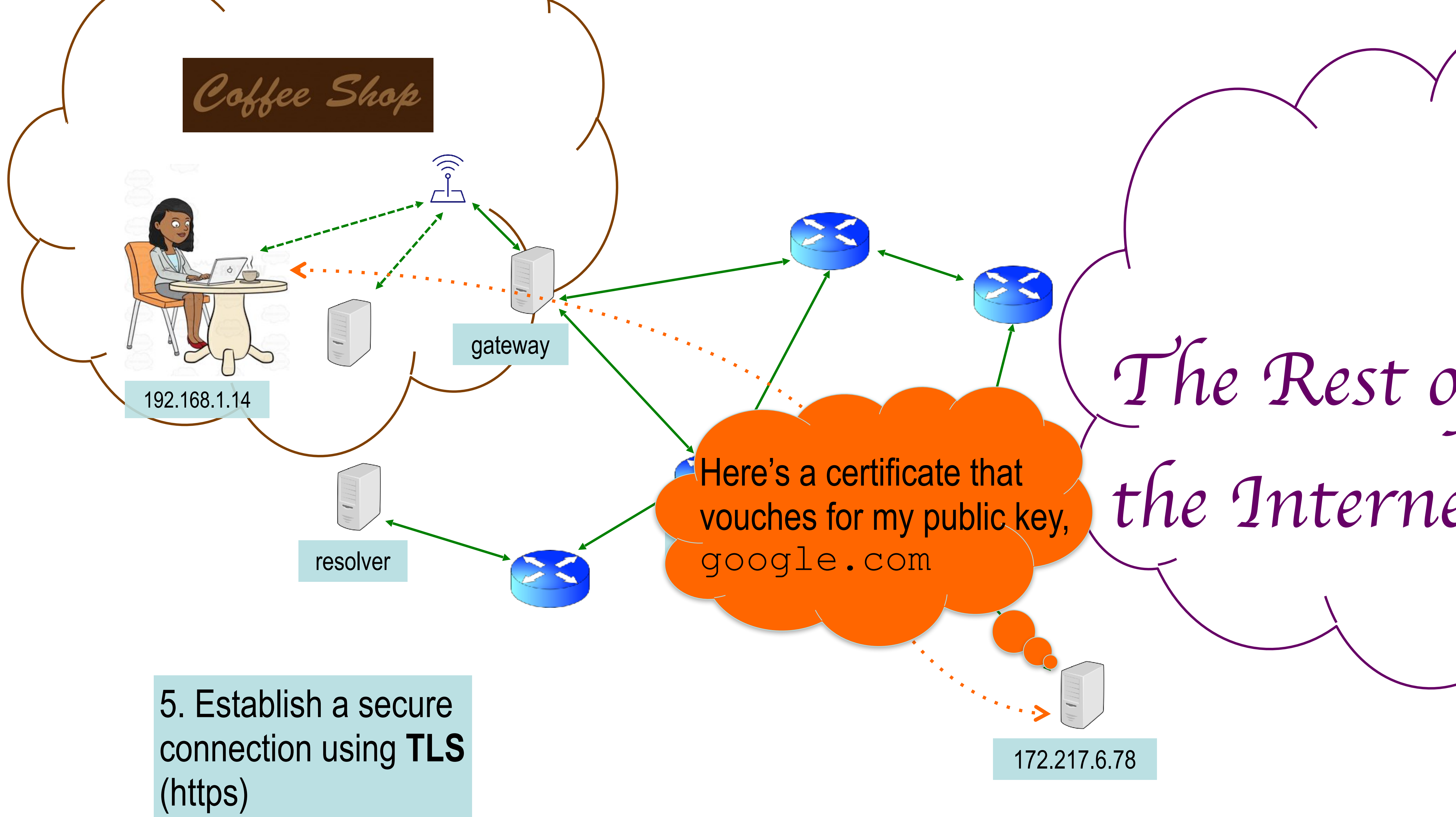
#### 4. Connect to google.com server



At this point the connection is established and data can be (reliably) exchanged.









Coffee Shop



192.168.1.14



gateway

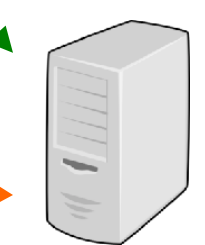
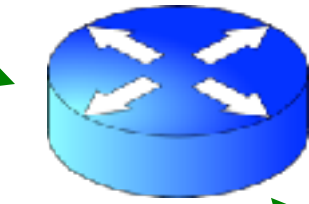


resolver

Well if you really possess the corresponding private key, prove it by decrypting this blob which we'll use to establish shared secret keys (RSA key exchange)



router



172.217.6.78

The Rest of the Internet

5. Establish a secure connection using **TLS** (https)

Coffee Shop

Well if you really possess the corresponding private key, prove it by signing your Diffie/Hellman half of a DHE using the private key?  
(DHE/ECDHE key exchange)



192.168.1.14



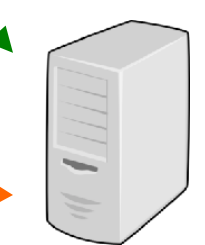
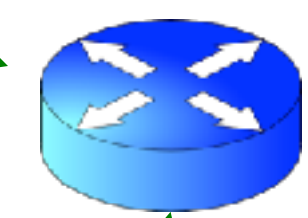
gateway



resolver



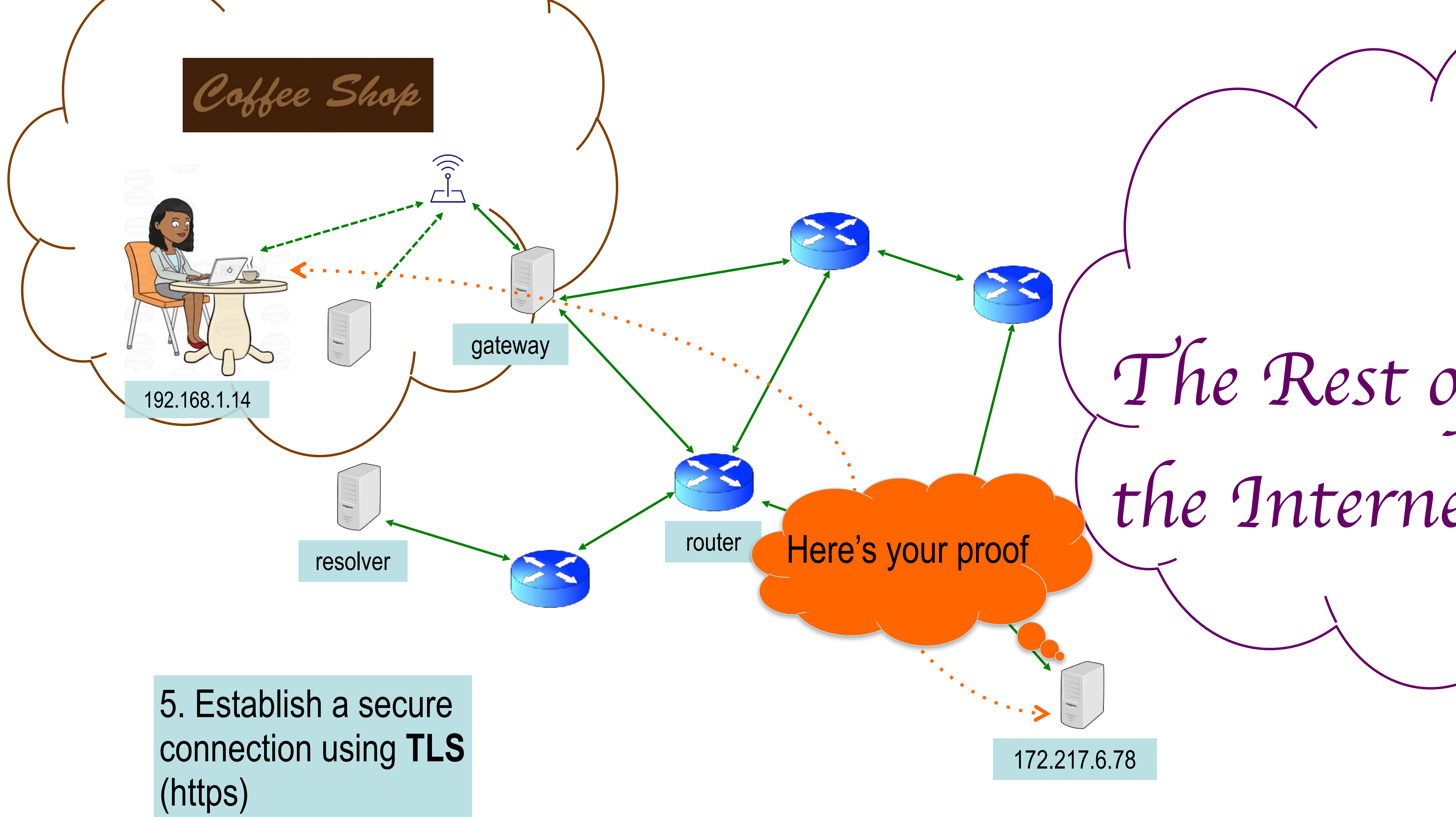
router



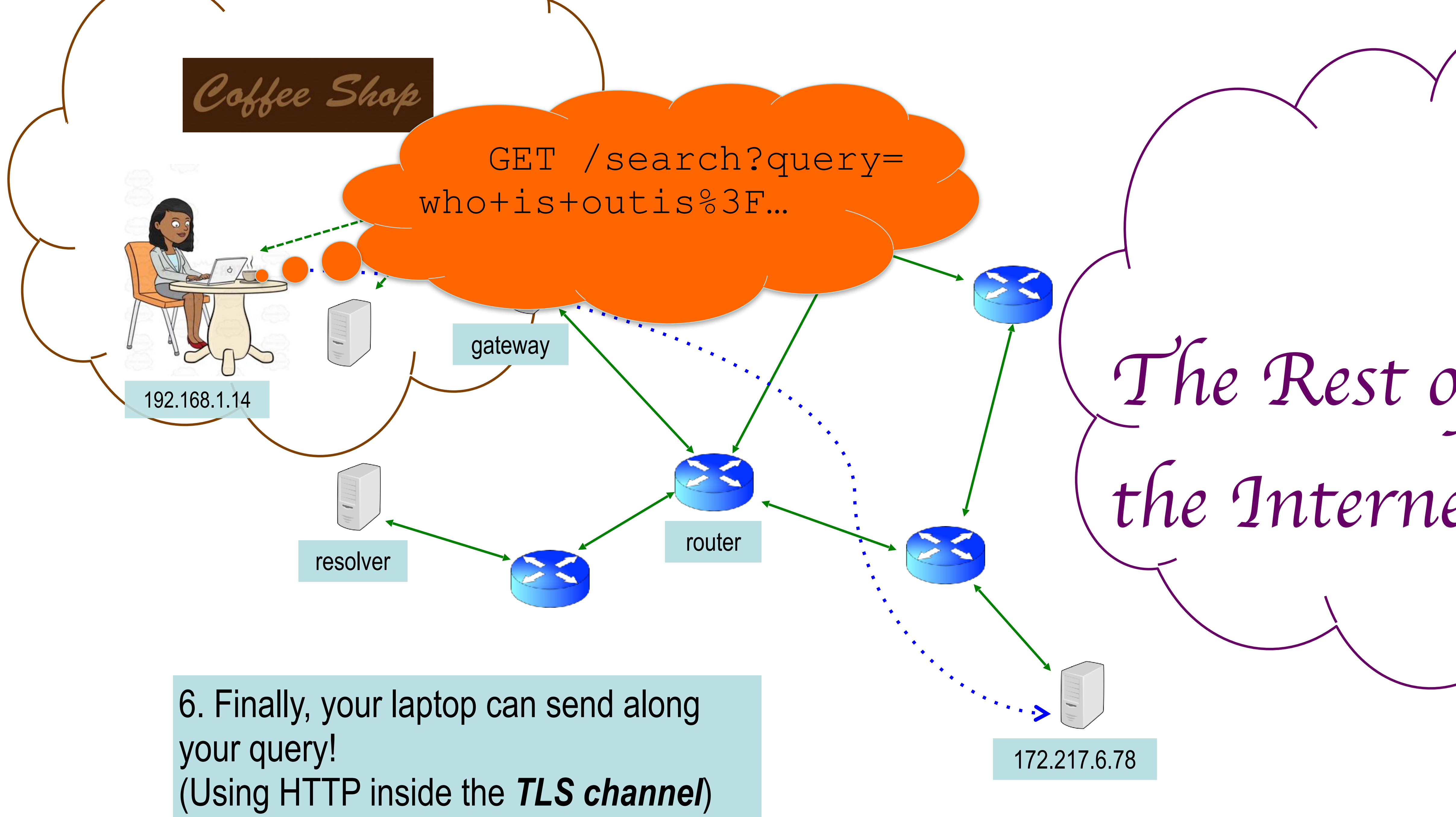
172.217.6.78

The Rest of the Internet

5. Establish a secure connection using **TLS** (https)







# IP addresses

- IPv4 addresses are 32 bits
  - **aa.bb.cc.dd**
    - Decimal values from 0-255, e.g. 128.32.131.12
- IPv6 addresses are 128 bits
  - **aaaa:bbbb:cccc:dddd:eeee:ffff:gggg:hhhh**
    - Hexadecimal values (can drop leading 0), e.g. 2607:f140:2000:4001:187f:86cc:3dfc:b9c8
    - A long run of 0s can be replaced with ::
- Subnets (/8, /16, /24...)
  - **128.32/16**
    - All IPv4 between 128.32.0.0 and 128.32.255.255
  - **2607:f140:2000:4001/64**
    - All IPv6 addresses with the same upper 64 bits



# Special IP addresses & Networks

- **Localhost: 127.0.0/24**
- **Broadcast: 255.255.255.255**
  - Send to all in the local network
  - Also for subnet, can specify all bits as 1 (e.g. for 128.32/16, 128.32.255.255) to broadcast to that network, but generally ignored these days
- **Private: 10/8, 172.16/12 (ends up being .16-.32), 192.168/16**
  - Not routed on the Internet, can use for internal purposes
  - Commonly used for NAT (more later)
- **IPv6 Multicast: ff00:/8**
  - In particular ff00::1 -> all machines on local network

# Ethernet

- 6 bytes of destination MAC address
  - In this case, MAC means media access control address, not message authentication code!
- 6 bytes of source MAC address
- Optional 4-byte VLAN tag
- 2-byte length/type field
- 46-1500 bytes of payload

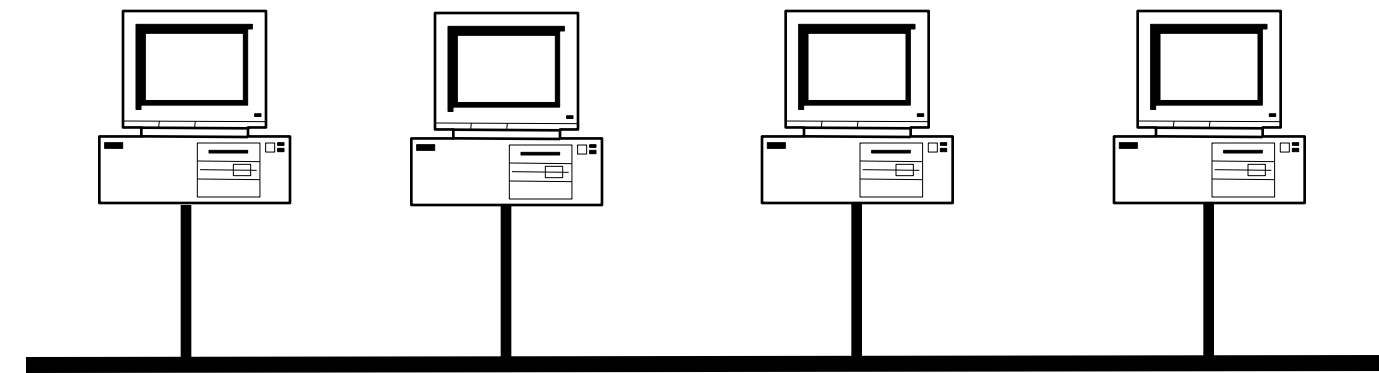
DST MAC	SRC MAC	VLAN	Type	PAYLOAD
---------	---------	------	------	---------

# The MAC Address

- The MAC address acts as a device identifier
  - Usually written as 6 bytes in hex, e.g. `13:37:ca:fe:f0:0d`
- A device ***should ignore*** all packets that aren't to itself or to the broadcast address (`ff:ff:ff:ff:ff:ff`)
  - But almost all devices can go into ***promiscuous mode***
    - This is also known as "sniffing traffic"
- A device generally should only send with its own address
  - But this is enforced with software and can be trivially bypassed when you need to write "raw packets"

# Attacks

# Link-layer threats



- Confidentiality: eavesdropping (aka sniffing)
- Integrity: injection of spoofed packets
- Availability: delete legit packets (e.g., jamming)



# Eavesdropping

- For subnets using broadcast technologies (e.g., WiFi, some types of Ethernet), attacker can eavesdrop
  - Each attached system's NIC (= Network Interface Card) can capture any communication on the subnet
  - Tools: tcpdump / windump (low-level text-based printout), wireshark (GUI)

# Wireshark

all.trace2 [Wireshark 1.6.2]						
File Edit View Go Capture Analyze Statistics Telephony Tools Internals Help						
Filter: Expression... Clear Apply						
No.	Time	Source	Destination	Protocol	Length	Info
1	0.000000	10.0.1.9	10.0.1.255	BJNP	58	Printer Command: Unknown code (2)
2	0.000198	10.0.1.9	224.0.0.1	BJNP	58	Printer Command: Unknown code (2)
3	2.150663	10.0.1.9	255.255.255.255	DB-LSP-D	172	Dropbox LAN sync Discovery Protocol
4	2.150938	10.0.1.9	10.0.1.255	DB-LSP-D	172	Dropbox LAN sync Discovery Protocol
5	4.514403	10.0.1.13	31.13.75.23	TCP	78	61901 > http [SYN] Seq=0 Win=65535 Len=0 MSS=1460 WS=8 TSval=4290
6	4.536771	31.13.75.23	10.0.1.13	TCP	74	http > 61901 [SYN, ACK] Seq=0 Ack=1 Win=14480 Len=0 MSS=1460 SACK
7	4.536896	10.0.1.13	31.13.75.23	TCP	66	61901 > http [ACK] Seq=1 Ack=1 Win=524280 Len=0 TSval=429017456 T
8	4.537429	10.0.1.13	31.13.75.23	HTTP	590	GET / HTTP/1.1
9	4.553947	31.13.75.23	10.0.1.13	TCP	66	http > 61901 [ACK] Seq=1 Ack=525 Win=15872 Len=0 TSval=1765827012
10	4.626447	31.13.75.23	10.0.1.13	HTTP	600	HTTP/1.1 302 Found
11	4.626579	10.0.1.13	31.13.75.23	TCP	66	61901 > http [ACK] Seq=525 Ack=535 Win=524280 Len=0 TSval=4290174
12	7.065664	10.0.1.9	10.0.1.255	BJNP	58	Printer Command: Unknown code (2)
13	7.065846	10.0.1.9	224.0.0.1	BJNP	58	Printer Command: Unknown code (2)

▶ Frame 10: 600 bytes on wire (4800 bits), 600 bytes captured (4800 bits)	
▶ Ethernet II, Src: Apple_fe:aa:41 (00:25:00:fe:aa:41), Dst: Apple_41:eb:00 (e4:ce:8f:41:eb:00)	
▶ Internet Protocol Version 4, Src: 31.13.75.23 (31.13.75.23), Dst: 10.0.1.13 (10.0.1.13)	
▼ Transmission Control Protocol, Src Port: http (80), Dst Port: 61901 (61901), Seq: 1, Ack: 525, Len: 534	
Source port: http (80)	
Destination port: 61901 (61901)	
[Stream index: 0]	
Sequence number: 1 (relative sequence number)	
[Next sequence number: 535 (relative sequence number)]	
Acknowledgement number: 525 (relative ack number)	
Header length: 32 bytes	
▶ Flags: 0x18 (PSH, ACK)	
Window size value: 31	
[Calculated window size: 15872]	
[Window size scaling factor: 512]	
▶ Checksum: 0xf42f [validation disabled]	
.....	
0000	e4 ce 8f 41 eb 00 00 25 00 fe aa 41 08 00 45 20 ...A...% ...A..E
0010	02 4a 67 be 00 00 58 06 83 9f 1f 0d 4b 17 0a 00 .Jg...X. ....K...
0020	01 0d 00 50 f1 cd d5 b8 c0 31 96 68 cb 28 80 18 ...P.... .l.h.(.
0030	00 1f f4 2f 00 00 01 01 08 0a 69 40 62 0b 19 92 .../.... .i@b...
0040	49 70 48 54 54 50 2f 31 2e 31 20 33 30 32 20 46 IpHTTP/1 .l 302 F
.....	
Frame (frame), 600 bytes	
Packets: 13 Displayed: 13 Marked: 0 Load time: 0:00.109	
Profile: Default	



## Operation Ivy Bells

*By Matthew Carle*  
**Military.com**

At the beginning of the 1970's, divers from the specially-equipped submarine, USS Halibut (SSN 587), left their decompression chamber to start a bold and dangerous mission, code named "Ivy Bells".



The Regulus guided missile submarine, USS Halibut (SSN 587) which carried out Operation Ivy Bells.



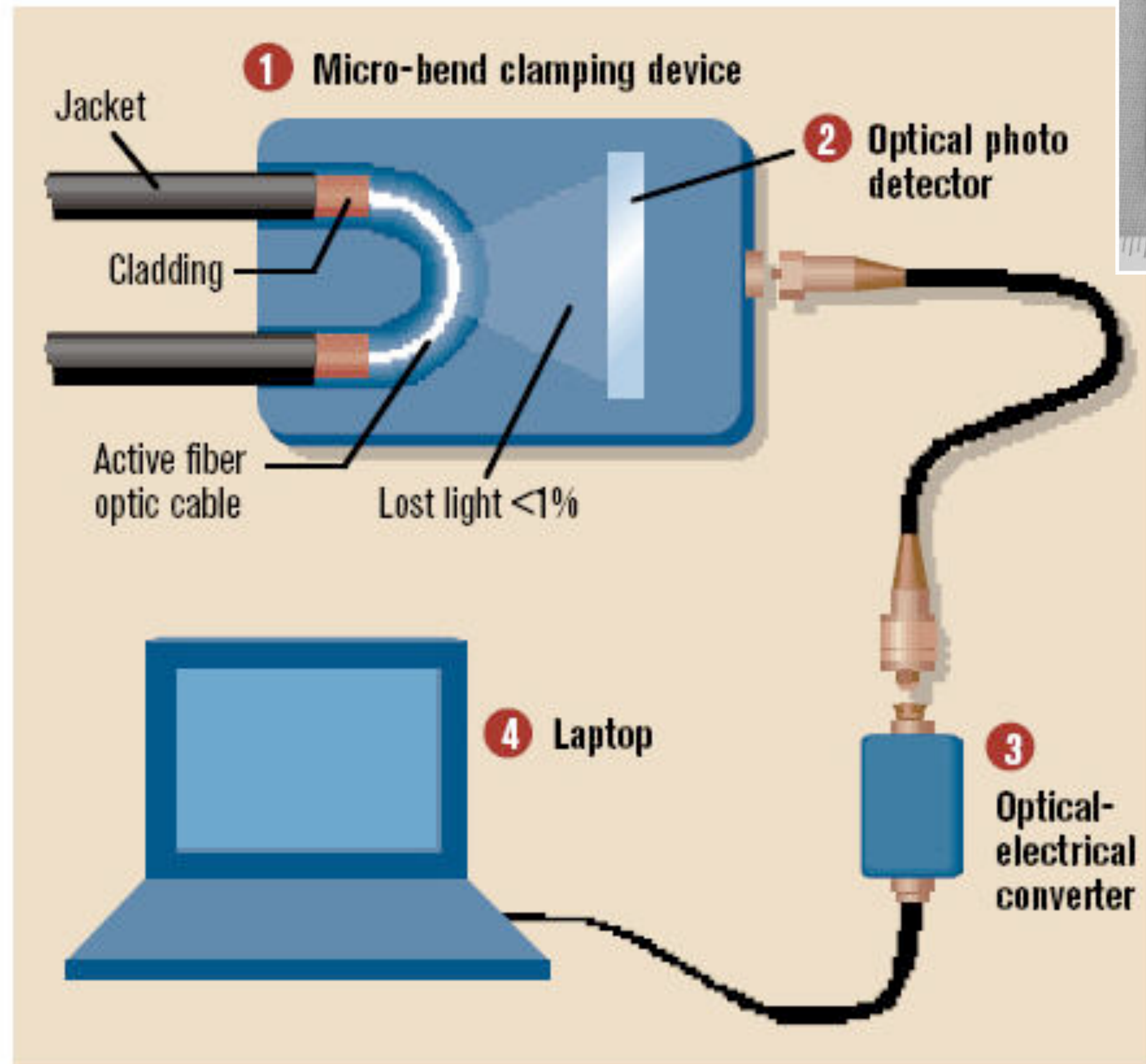
In an effort to alter the balance of Cold War, these men scoured the ocean floor for a five-inch diameter cable carry secret Soviet communications between military bases.

The divers found the cable and installed a 20-foot long listening device on the cable. designed to attach to the cable without piercing the casing, the device recorded all communications that occurred. If the cable malfunctioned and the Soviets raised it for repair, the bug, by design, would fall to the bottom of the ocean. Each month Navy divers retrieved the recordings and installed a new set of tapes.

Upon their return to the United States, intelligence agents from the NSA analyzed the recordings and tried to decipher any encrypted information. The Soviets apparently were confident in the security of their communications lines, as a surprising amount of sensitive information traveled through the lines without encryption.

prison. The original tap that was discovered by the Soviets is now on exhibit at the KGB museum in Moscow.







# Link-Layer Spoofing

- Attacker can inject spoofed packets, and lie about the source address

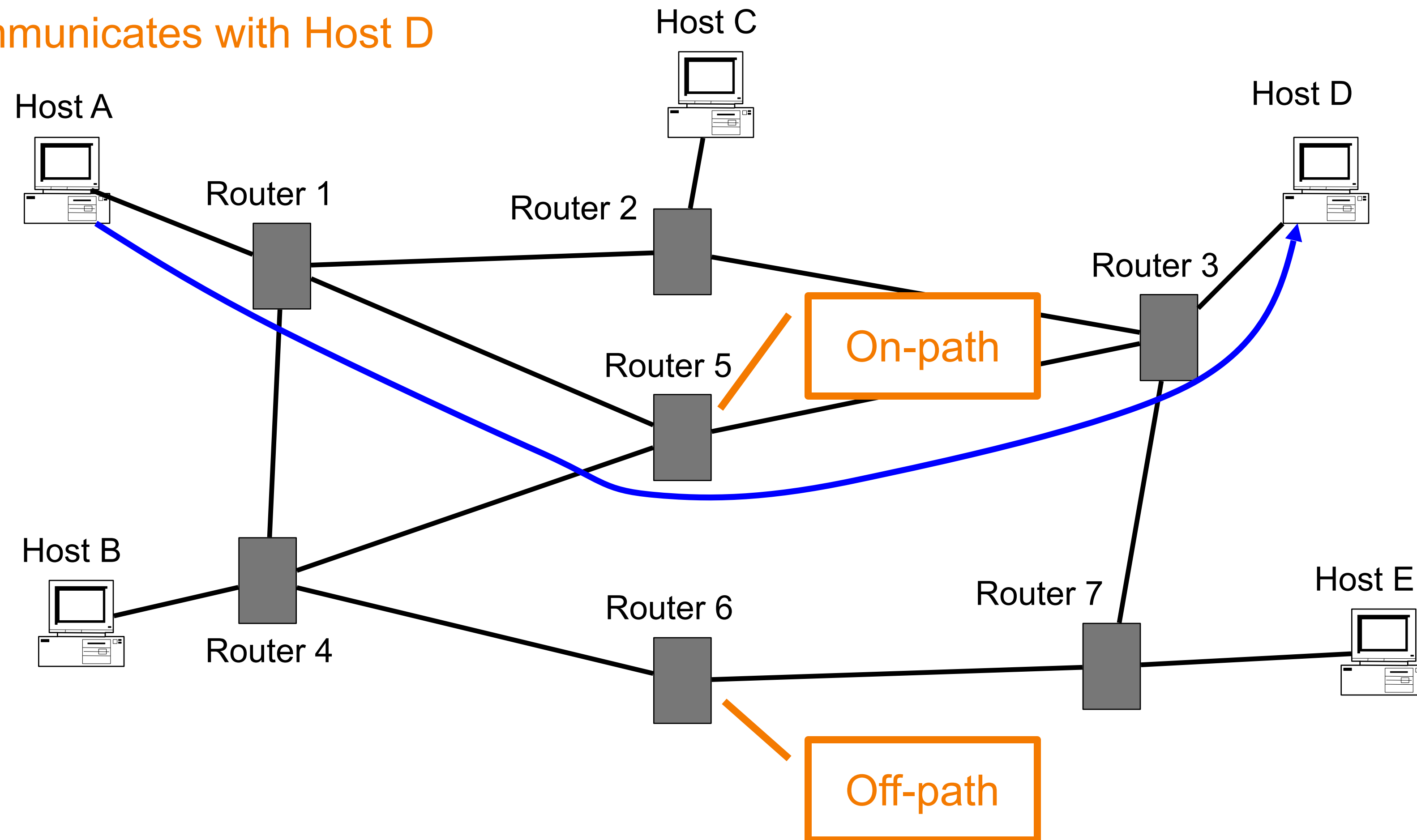


# Spoofing

- With physical access to a local network, attacker can create any packet they like
  - Spoofing = lie about source address
- Particularly powerful when combined with eavesdropping, because attacker can understand exact state of victim's communication and craft their spoofed traffic to match it
  - Spoofing w/o eavesdropping = blind spoofing

# On-path vs Off-path Spoofing

Host A communicates with Host D



# Spoofing on the Internet

- On-path attackers can see victim's traffic  $\Rightarrow$  spoofing is easy
- Off-path attackers can't see victim's traffic
  - They have to resort to blind spoofing
  - Often must guess/infer header values to succeed
    - We then care about work factor: how hard is this
  - But sometimes they can just brute force
    - E.g., 16-bit value: just try all 65,536 possibilities!
- When we say an attacker “can spoof”, we usually mean “w/ reasonable chance of success”



# DNS Service

- Runs Domain Name Servers
- Translates domain names google.com to IP addresses
- When user browser wants to contact google.com, it first contacts a DNS to find out the IP address for google.com and then sends a packet to that IP address
- More in future lectures..